



NTNU
Norwegian University of
Science and Technology

Towards Privacy Preserving Mobile Communications

George Petrides

(with Kristian Gjøsteen and Asgeir Steine)

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Information Security Research School
Finse

Current Situation

(very loosely speaking)

Mobile Network Operators (MNOs) – e.g. Telenor, Netcom:

- 1 Maintain the communication infrastructure (base stations).
- 2 Subscribe users (SIM cards with embedded symmetric key & IMSI).
- 3 Bill subscribers and other MNOs for services.
 - Virtual Network Operators (MVNOs) – e.g. Chess.

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 - Eavesdroppers can't follow Alice around.
 - \mathcal{N}_A can!
 - Active attackers can too! (IMSI-catchers)

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If Alice wants to call Bob:

- Alice updates her position to \mathcal{N}_A and asks to contact Bob -
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\mathcal{N}_A and \mathcal{N}_B (possibly one and the same) learn EVERYTHING there is to know: Who, Where and What.

New Privacy Preserving Setting

- We would prefer if \mathcal{N}_A can't follow Alice around
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Proposal: Split MNOs in two: MNOs and SPs (Service Providers)

1 MNOs

- Maintain the communication infrastructure (base stations).
- Bill SPs for services.

2 SPs

- Subscribe users (SIM card with embedded symmetric key & identity token).
- Bill subscribers for services.

This is not crazy - similar to MNO–MVNO case!

User-MNO Key Establishment

Alice establishes a secure channel with nearest \mathcal{N} .

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- 1 Alice (user ps_A) identifies herself to \mathcal{SP}_A using token

$$\mathcal{T}_A = \text{Enc}\{Alice || smth\}_{\mathcal{K}_{\mathcal{SP}_A}}.$$

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\mathcal{N} - \mathcal{SP}_A collusion leaks all!

User-MNO Key Establishment - The Protocol

Alice
(ps_A)

\mathcal{N}

SP_A

(T_A, n_A, g^x, SP_A)
→

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Alice
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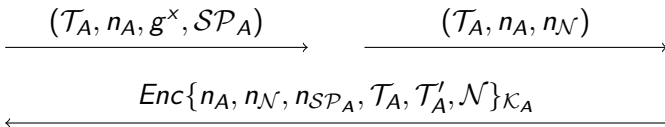


User-MNO Key Establishment - The Protocol

Alice
(p_{SA})

\mathcal{N}

SP_A

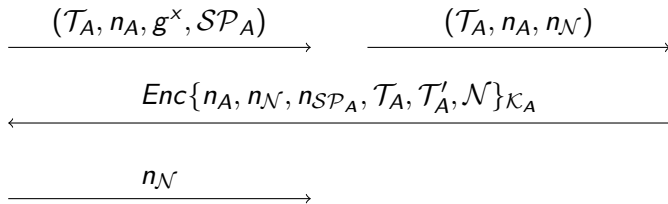


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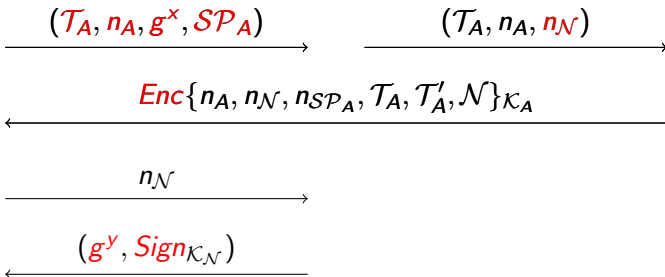


User-MNO Key Establishment - The Protocol

Alice
(p_{SA})

\mathcal{N} 😊

SP_A

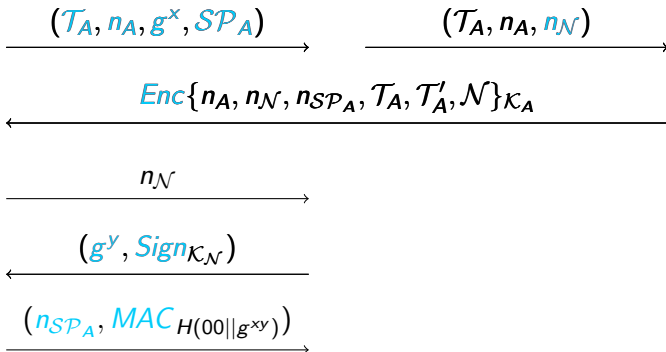


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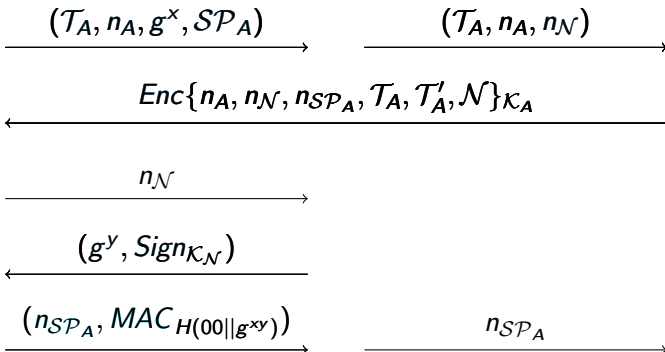


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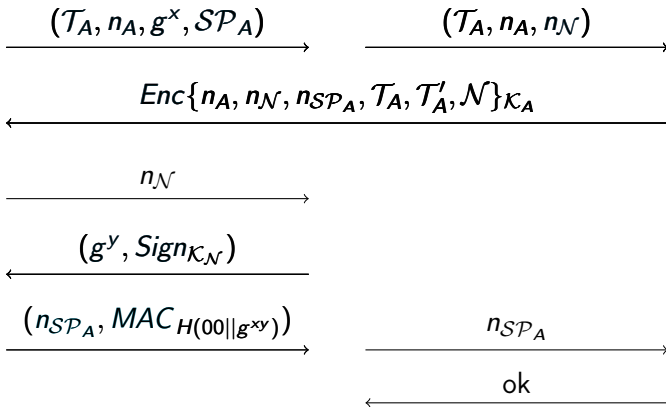


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SP_A 😊

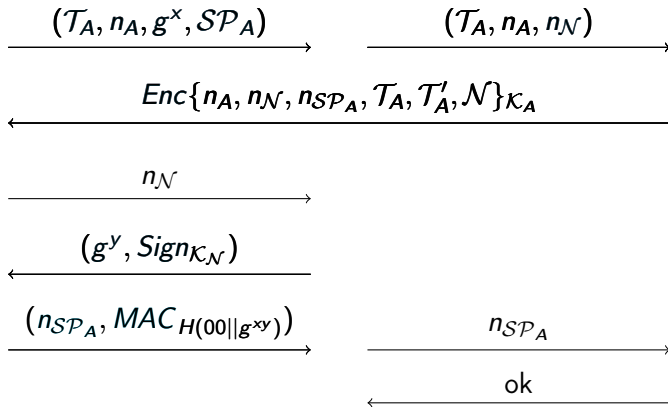


User-MNO Key Establishment - The Protocol

Alice 😊
(p_{SA})

\mathcal{N} 😊

SP_A 😊



Shared: keys $H(01||g^{xy})$, $H(10||g^{xy})$, TMSI $H(11||g^{xy})$

Traceability Issues

Alice uses token \mathcal{T}_A to verify that she is a subscriber.

- To avoid tracing, \mathcal{SP}_A sends her a new token \mathcal{T}_A' .
- Still, traceable if denial of service before she receives \mathcal{T}_A' .
 - Only if she moves before connecting.
- Unconditional untraceability alternative: Public Key Crypto.
 - Too expensive to verify valid tokens - Denial of service attack!

Authenticated & Encrypted Radio Link

- Using the established keys and TMSI Alice secures the radio link with \mathcal{N} .
- \mathcal{N} sends pseudonyms to Alice for communication with others.
- Alice can have persistent connection with content service providers.
 - Initial authentication with PKC (no DoS attacks as users are tied to TMSI via pseudonym).
 - Stay connected using tokens (if new token denied, go PKC again).
 - Protection against traffic analysis.
- Can contact Bob by requesting his pseudonym from his Telephony (content) Provider.
 - Again, initially PKC and then tokens.
 - No one knows who is calling who.
 - Can encrypt their communication using a symmetric key.

Legal Aspects

What about e.g. EU's Data Retention Directive?

- Judicial MNO–SP collaboration can reveal necessary info.
- Stored info is split and therefore leaks less about users.

Conclusions

- Don't expect implementation - mainly make a point.
- Designed to be Universally Composable Secure
 - We provide ideal functionalities for everything.
 - Proofs are quite long and complex (as the case usually is).

Thank You!