

Prototype of Russian Hash Function "Stribog"

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Outline

- 1 Introduction
- 2 Description of Stribog
- 3 Performance Comparison With Other Hash Functions
- 4 Conclusions

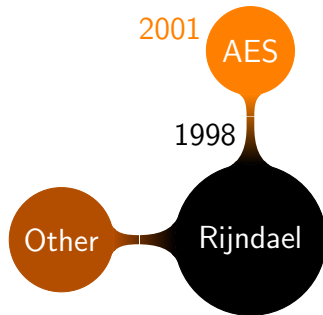
- GOST 34.11-94 was **theoretically broken** in 2008.
 - The complexities $O(2^{192})/O(2^{69})$ for preimage and second preimage attacks.

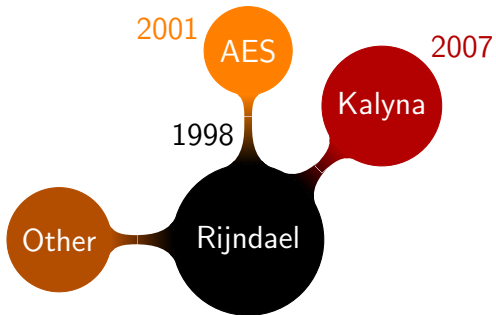
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- **Increasing performance.** Stribog is 20% faster than GOST 34.11-94.

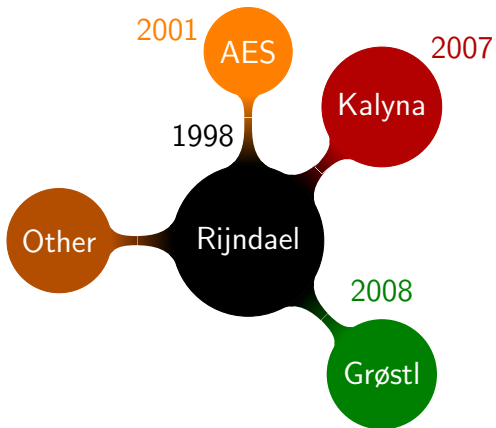
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- Opposite to SHA-3.

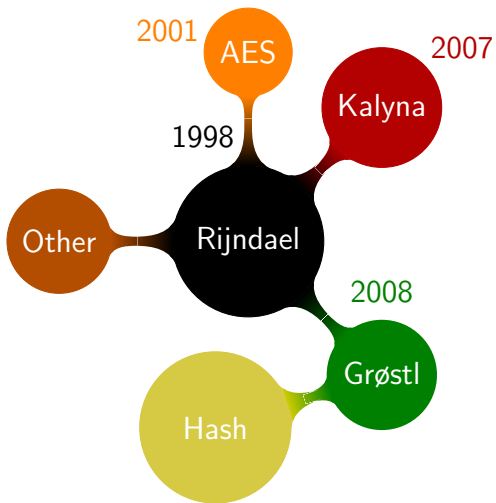
1998

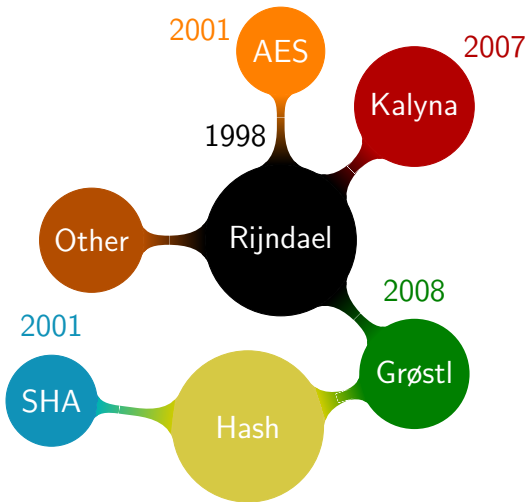


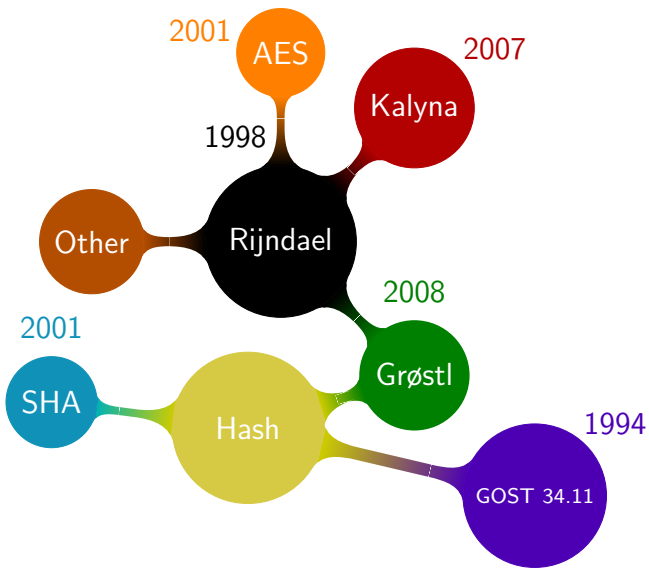


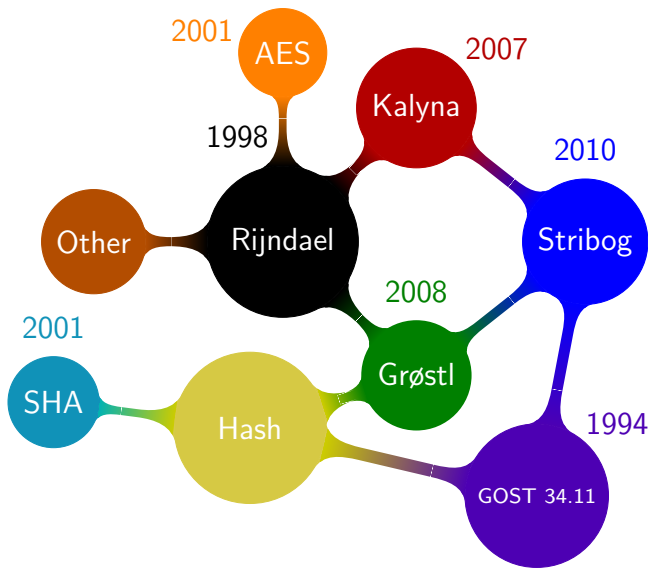












Basic Operations and Functions

Stribog is based on SP-network block cipher with block and key length equal 512 bits

- **SubBytes (S)**: nonlinear bijective mapping.
- **Transposition (P)**: byte permutation.
- **MixColumns (L)**: linear transformation.
- **AddRoundKey (X)**: addition with the round key using bitwise XOR.

Other basic functions

- \boxplus : addition modulo 2^{512} .
- $MSB_s(A)$: getting s most significant bits of vector A .
- $A||B$: concatenation of two vectors A and B .

State Representation

Grøstl

a_0	a_8	a_{16}	a_{24}	a_{32}	a_{40}	a_{48}	a_{56}
a_1	a_9	a_{17}	a_{25}	a_{33}	a_{41}	a_{49}	a_{57}
a_2	a_{10}	a_{18}	a_{26}	a_{34}	a_{42}	a_{50}	a_{58}
a_3	a_{11}	a_{19}	a_{27}	a_{35}	a_{43}	a_{51}	a_{59}
a_4	a_{12}	a_{20}	a_{28}	a_{36}	a_{44}	a_{52}	a_{60}
a_5	a_{13}	a_{21}	a_{29}	a_{37}	a_{45}	a_{53}	a_{61}
a_6	a_{14}	a_{22}	a_{30}	a_{38}	a_{46}	a_{54}	a_{62}
a_7	a_{15}	a_{23}	a_{31}	a_{39}	a_{47}	a_{55}	a_{63}



$$A = a_0 || a_1 || \dots || a_{63}$$

Stribog

b_{63}	b_{62}	b_{61}	b_{60}	b_{59}	b_{58}	b_{57}	b_{56}
b_{55}	b_{54}	b_{53}	b_{52}	b_{51}	b_{50}	b_{49}	b_{48}
b_{47}	b_{46}	b_{45}	b_{44}	b_{43}	b_{42}	b_{41}	b_{40}
b_{39}	b_{38}	b_{37}	b_{36}	b_{35}	b_{34}	b_{33}	b_{32}
b_{31}	b_{30}	b_{29}	b_{28}	b_{27}	b_{26}	b_{25}	b_{24}
b_{23}	b_{22}	b_{21}	b_{20}	b_{19}	b_{18}	b_{17}	b_{16}
b_{15}	b_{14}	b_{13}	b_{12}	b_{11}	b_{10}	b_9	b_8
b_7	b_6	b_5	b_4	b_3	b_2	b_1	b_0



$$B = b_{63} || b_{62} || \dots || b_0$$

State Representation

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a_1	a_9	a_{17}	a_{25}	a_{33}	a_{41}	a_{49}	a_{57}
a_2	a_{10}	a_{18}	a_{26}	a_{34}	a_{42}	a_{50}	a_{58}
a_3	a_{11}	a_{19}	a_{27}	a_{35}	a_{43}	a_{51}	a_{59}
a_4	a_{12}	a_{20}	a_{28}	a_{36}	a_{44}	a_{52}	a_{60}
a_5	a_{13}	a_{21}	a_{29}	a_{37}	a_{45}	a_{53}	a_{61}
a_6	a_{14}	a_{22}	a_{30}	a_{38}	a_{46}	a_{54}	a_{62}
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Stribog

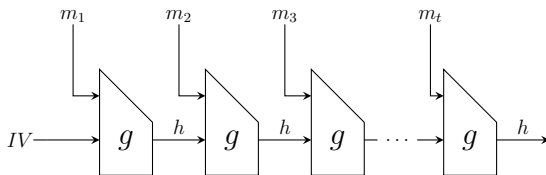
a_0	a_1	a_2	a_3	a_4	a_5	a_6	a_7
a_8	a_9	a_{10}	a_{11}	a_{12}	a_{13}	a_{14}	a_{15}
a_{16}	a_{17}	a_{18}	a_{19}	a_{20}	a_{21}	a_{22}	a_{23}
a_{24}	a_{25}	a_{26}	a_{27}	a_{28}	a_{29}	a_{30}	a_{31}
a_{32}	a_{33}	a_{34}	a_{35}	a_{36}	a_{37}	a_{38}	a_{39}
a_{40}	a_{41}	a_{42}	a_{43}	a_{44}	a_{45}	a_{46}	a_{47}
a_{48}	a_{49}	a_{50}	a_{51}	a_{52}	a_{53}	a_{54}	a_{55}
a_{56}	a_{57}	a_{58}	a_{59}	a_{60}	a_{61}	a_{62}	a_{63}

$$A = a_0 || a_1 || \dots || a_{63}$$

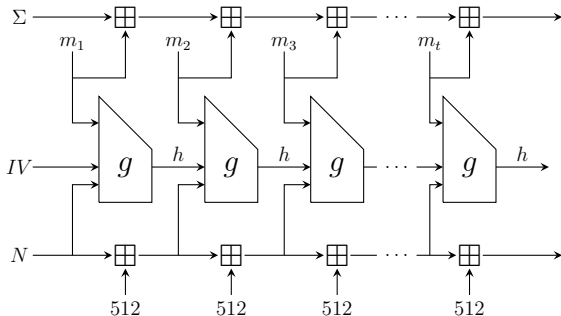
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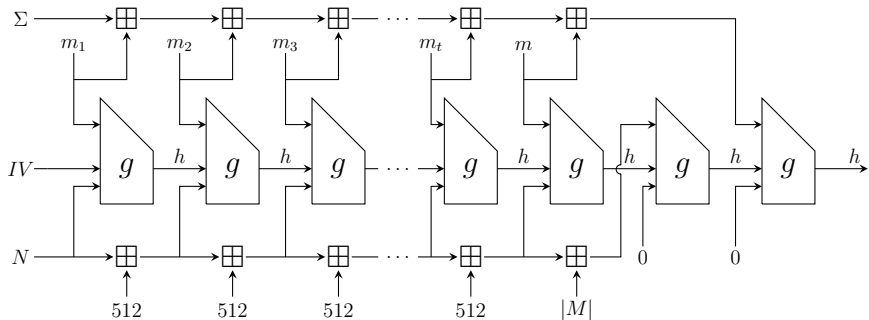
Merkle-Damgård Scheme



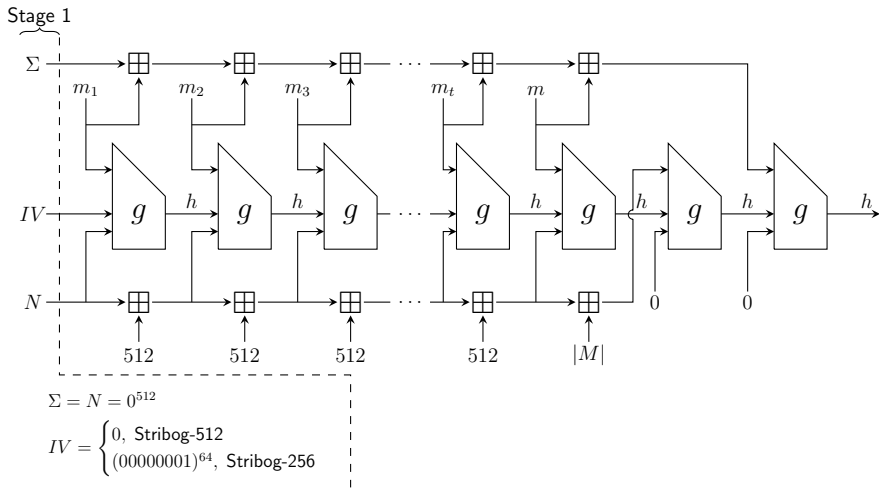
Modification of Merkle-Damgård Scheme



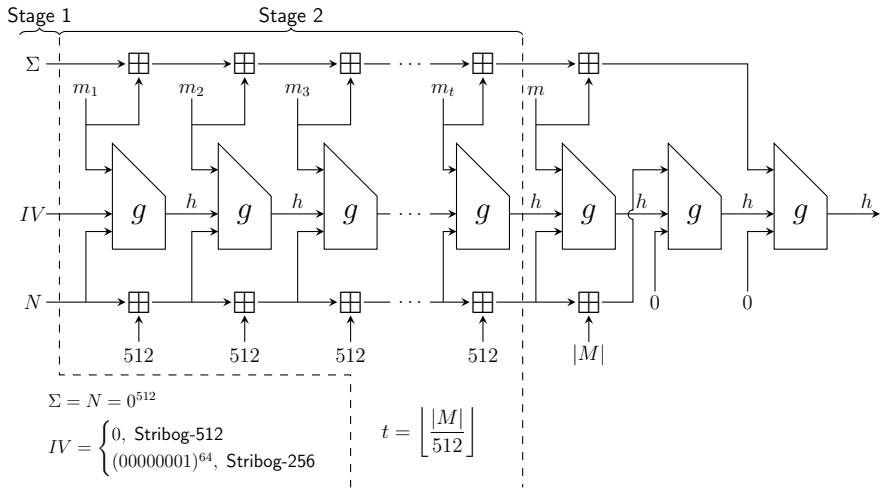
Hash Function Stribog



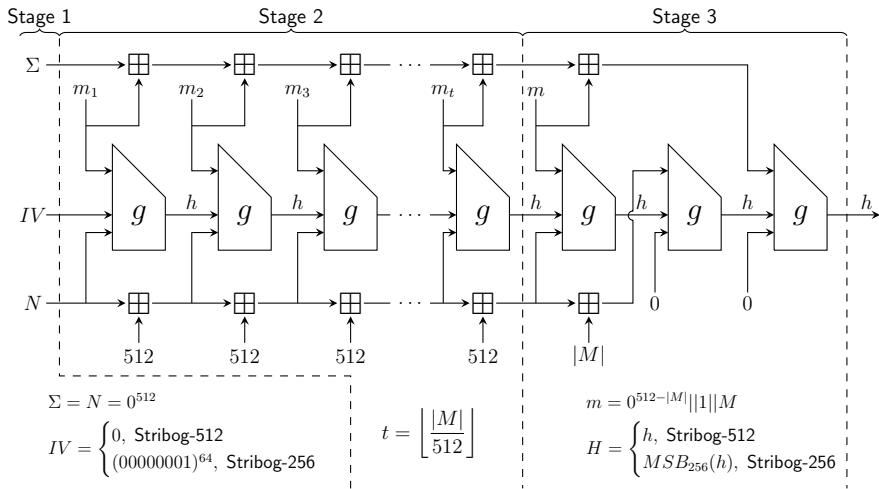
Hash Function Stribog. Stage 1



Hash Function Stribog. Stage 2

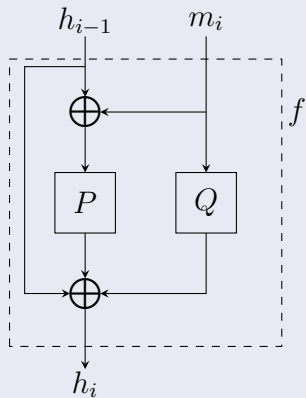


Hash Function Stribog. Stage 3

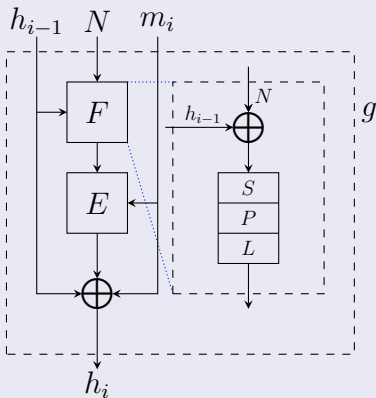


Compression Function Construction

Grøstl



Stribog



Design of E

Compression function $g_N : \mathbb{F}_2^{512} \times \mathbb{F}_2^{512} \mapsto \mathbb{F}_2^{512}$, $N \in \mathbb{F}_2^{512}$ is defined as follows

$$g_N(h, m) = E(L \circ P \circ S(h \oplus N), m) \oplus h \oplus m, \quad h, m \in \mathbb{F}_2^{512}$$

where

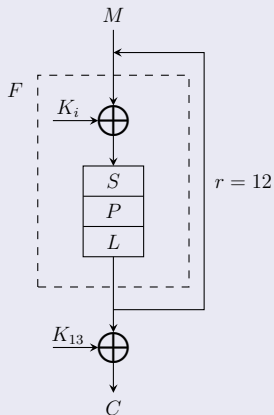
$$E(K, m) = X[K_{13}] \circ \prod_{i=1}^{12} L \circ P \circ S \circ X[K_i]$$

KeySchedule function

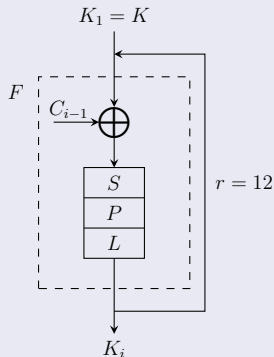
$$K_i = L \circ P \circ S(K_{i-1} \oplus C_{i-1}), \quad K_1 = K, \quad i \in \{2, \dots, 13\}.$$

Representation of E

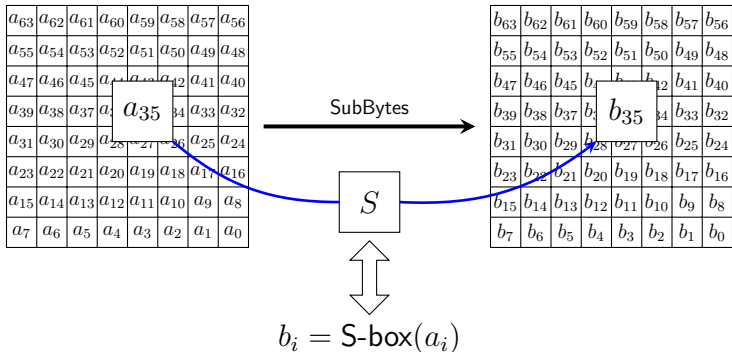
StribogT BC



Key Schedule



SubBytes Transformation



S-box of Stribog

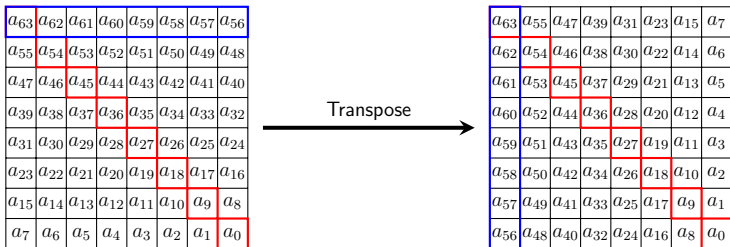
	0	1	2	3	4	5	6	7	8	9	A	B	C	D	E	F
0	FC	EE	DD	11	CF	6E	31	16	FB	C4	FA	DA	23	C5	04	4D
1	E9	77	F0	DB	93	2E	99	BA	17	36	F1	BB	14	CD	5F	C1
2	F9	18	65	5A	E2	5C	EF	21	81	1C	3C	42	8B	01	8E	4F
3	05	84	02	AE	E3	6A	8F	A0	06	0B	ED	98	7F	D4	D3	1F
4	EB	34	2C	51	EA	C8	48	AB	F2	2A	68	A2	FD	3A	CE	CC
5	B5	70	0E	56	08	0C	76	12	BF	72	13	47	9C	B7	5D	87
6	15	A1	96	29	10	7B	9A	C7	F3	91	78	6F	9D	9E	B2	B1
7	32	75	19	3D	FF	35	8A	7E	6D	54	C6	80	C3	BD	0D	57
8	DF	F5	24	A9	3E	A8	43	C9	D7	79	D6	F6	7C	22	B9	03
9	E0	0F	EC	DE	7A	94	B0	BC	DC	E8	28	50	4E	33	0A	4A
A	A7	97	60	73	1E	00	62	44	1A	B8	38	82	64	9F	26	41
B	AD	45	46	92	27	5E	55	2F	8C	A3	A5	7D	69	D5	95	3B
C	07	58	B3	40	86	AC	1D	F7	30	37	6B	E4	88	D9	E7	89
D	E1	1B	83	49	4C	3F	F8	FE	8D	53	AA	90	CA	D8	85	61
E	20	71	67	A4	2D	2B	09	5B	CB	9B	25	D0	BE	E5	6C	52
F	59	A6	74	D2	E6	F4	B4	C0	D1	66	AF	C2	39	4B	63	B6

S-box Characteristics

Characteristics	Stribog	AES
Boolean Functions		
Balanced	True	True
Nonlinearity	100	112
AI	96	32
SSI	258688	133120
PC	0	0
CI	0	0
Degree	7	7
Resiliency	0	0
SAC	False	False
Substitution		
Bijection	True	True
MDT	8	4
MLT	28	16
Cycles	252:243, 46:13	43:27, 242:87, 99:59, 124:81, 143:2
AI	3(441)	2(39)

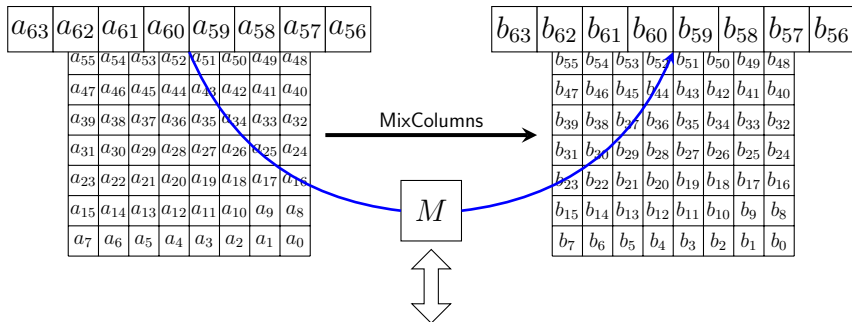
Transposition

Transposition transformation has a form



MixColumns

MixColumns transformation has a form



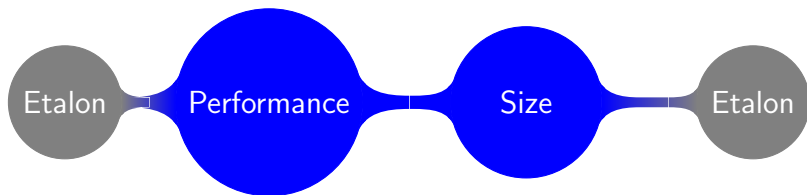
Multiplying the vector by the constant 64×64 matrix M in \mathbb{F}_2

$$B = A \cdot M$$

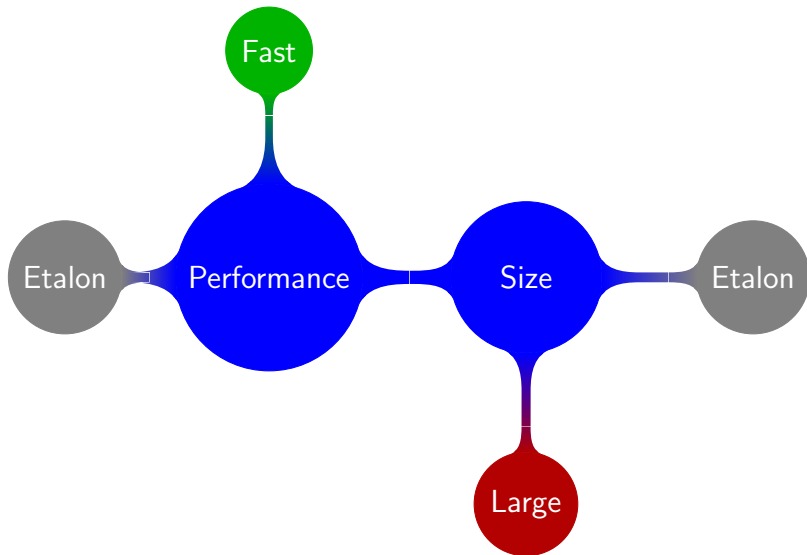
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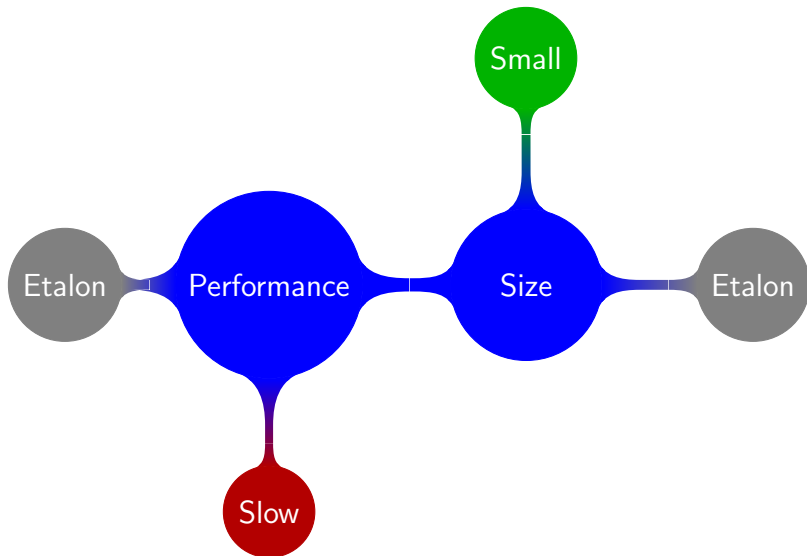
Implementation Features



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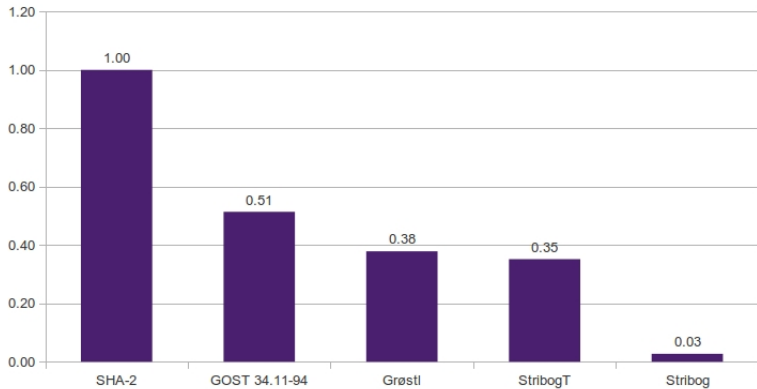
Implementation Features



Algorithm	Performance Mbits/s
SHA-2	74
GOST 34.11-94	38
Grøstl	28
StribogT	26
Stribog	2

Stribog with table implementation is 30% slower than GOST 34.11-94.

Performance



Comments for Performance

	Stribog BC	AES
Performance	140 Mbits/s	490 Mbits/s
Number of F	12	10
Number of lookup tables	64	16
Number of lookup tables (total)	768	160
Optimized	64 bit	32 bit

CPU: Intel T2050 1.60GHz **RAM:** 2GB

OS: Linux i686 **GCC:** version 4.6.3

Conclusions

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 - No, it is 30% slower on x86_32.
- Several mistakes in appendix (test vectors).
 - Appendices are for reference purposes only and are not part of the standard.

What is Stribog?



"Stribog in the Slavic pantheon, is the god and spirit of the winds, sky and air; he is said to be the ancestor (grandfather) of the winds of the eight directions."

– Wikipedia

References

-  F. Mendel, N. Pramstaller, C. Rechberger, M. Kontak, and J. Szmids. Cryptanalysis of the GOST hash function. In D. Wagner, editor, *Advances in Cryptology CRYPTO 2008*, volume 5157 of *LNCS*, pages 162–178.
-  Matuhin D.V., Shyshkin V.A., Rudskoy V.I.: Prospective hashing algorithm. RusCrypto'2010, 2010. (In Russian).
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