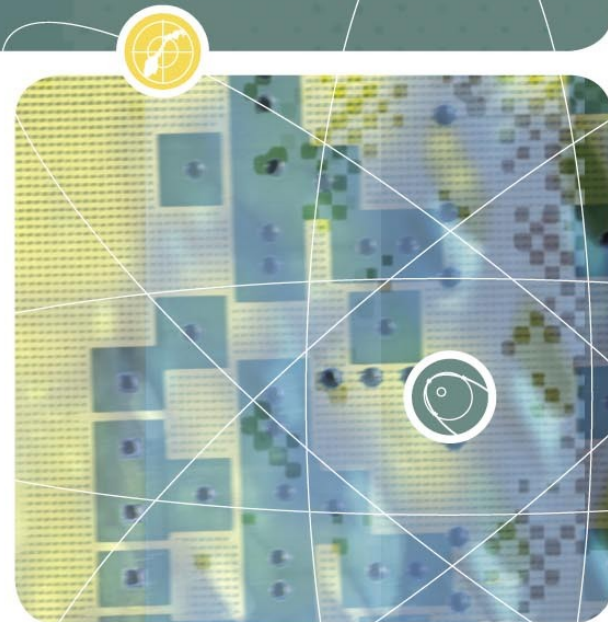


Identity Management Systems for multi-domain and tactical networks

Finseskolen 2012

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Norwegian Defence
Research
Establishment
May 2012



Why Identity Management (IdM)?

- Because we need integrated management of info about
 - Users
 - Equipment
 - Roles and Privileges
- A PKI (Public Key Infrastructure) only manages keys
- **IdM = Identity Provision+Authentication+Access Control**

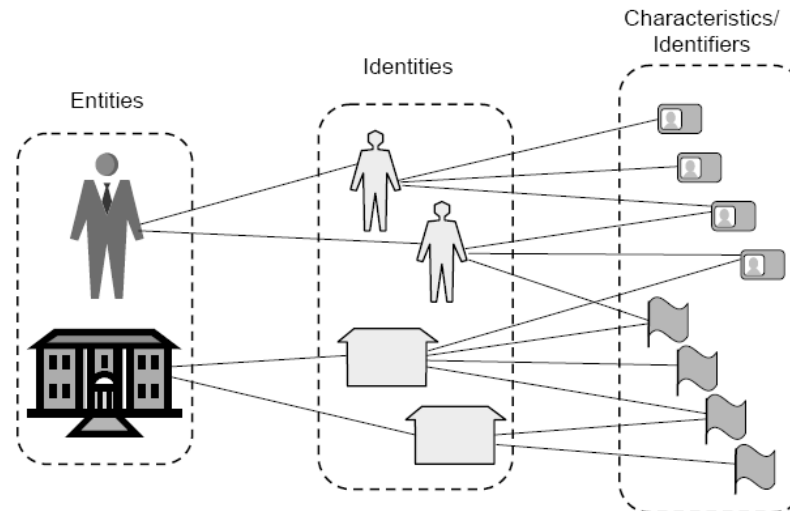


Figure 1: Correspondence between entities, identities and characteristics/identifiers.

Identity Management (IdM)

- Identity:
 - Set of properties associated with an *Entity*
- Identifier:
 - Subset of properties to *distinguish* identities
- Identity Statement:
 - Attestation of the subject's identifier
- Identity Provider (IdP)
 - Service which issues identity statements
- Authentication
 - Establishment of identity

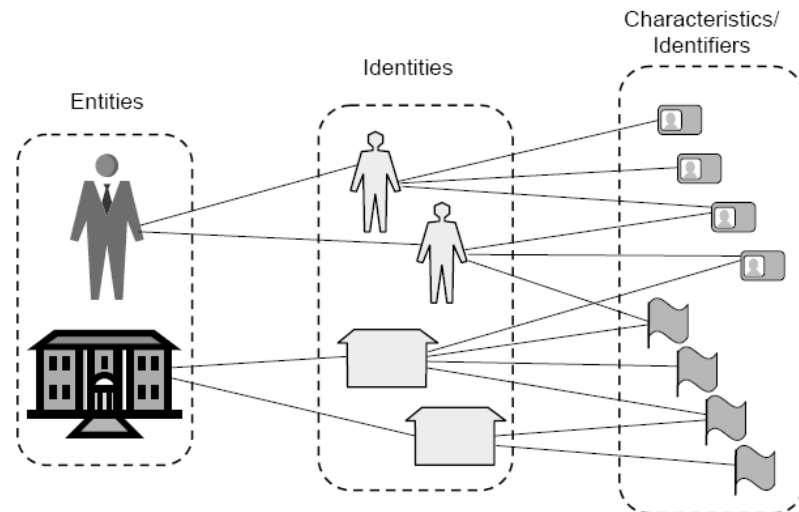


Figure 1: Correspondence between entities, identities and characteristics/identifiers.



IdM for mobile/tactical systems

- Are likely targets for intrusion attacks
- Highly dynamic client population
 - Even coalition networks of “guest” client
- Lightweight client units
 - Portability issues
 - Library availability
 - SOAP support
- Communication network
 - Availability/connectivity
 - Latency (round-trip delay)
 - Bandwidth (communication delay)



IdMs are suffering from:

- Discarding existing investments
 - need separate user registries
- High coupling between domains
 - guest users individually registered
 - autonomy delegated for federation
- Visibility of user identities
 - access given to identities, not roles
- Driven by security excellence, not networking excellence
 - protocols too costly for "narrow and bumpy" networks

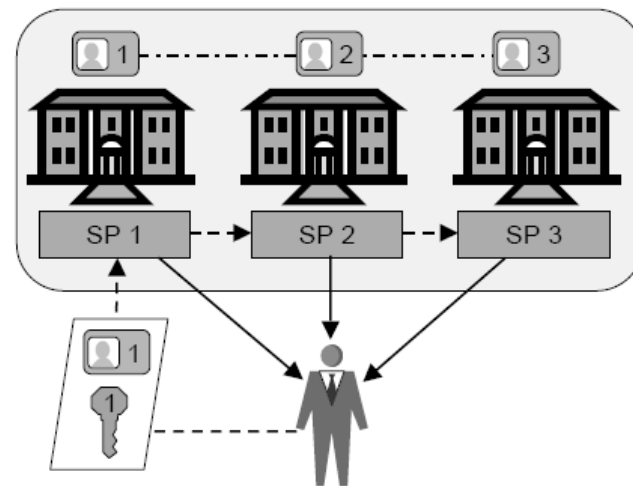


Figure 3: Federated user identity model.



IdM systems should

- Protect investments and knowledge
 - Employ existing enrollment procedures and data storage
- Allow federation for "guest access"
 - Should not need to enroll guests
- Give access rights to *roles*, not identities
 - RBAC, ABAC
- Protect domain autonomy
 - owner of service decides the access control
- Allow system latency
 - trust has a lifetime
- Limit the trust relationships
 - minimize the "trust anchors"
- Balance requirements between security and network economy



Tactical networks – ubiquitous computing

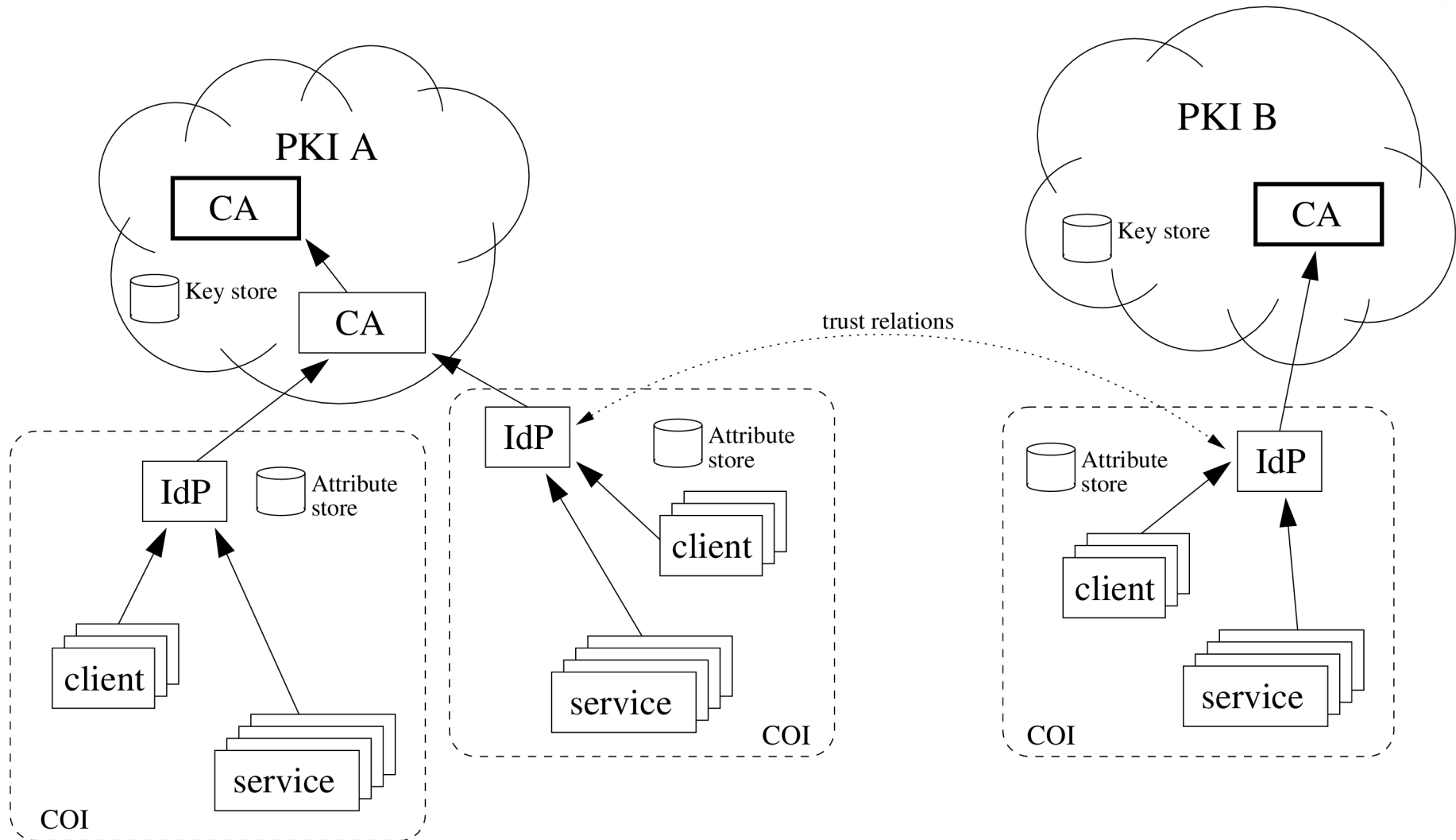
- Mobile, wireless, based on military radio technology
 - spread spectrum, strong encryption
- Low bandwidth (< 100 kb/s, depending on range)
- Multi-hop, Ad-hoc
 - latency
 - packet loss
 - link loss
- Applications adapted for tactical networks are frugal, robust and perserverant, which are desirable properties everywhere
 - ***tactical applications are fit for ubiquitous computing***



Revocation and Tactical Networks

- Identity credentials may need to be *revoked*
- Revocation of identity information requires bandwidth and connectivity
- Revocation checking is expensive and error-prone
 - since one actually asks the *opposite question*
- The work presented
 - relies on short-lived ***"identity statements"*** which require no revocation scheme,
 - the identity statements are derived from X.509 certificates maintained in a PKI

The GISMO IdM Architecture





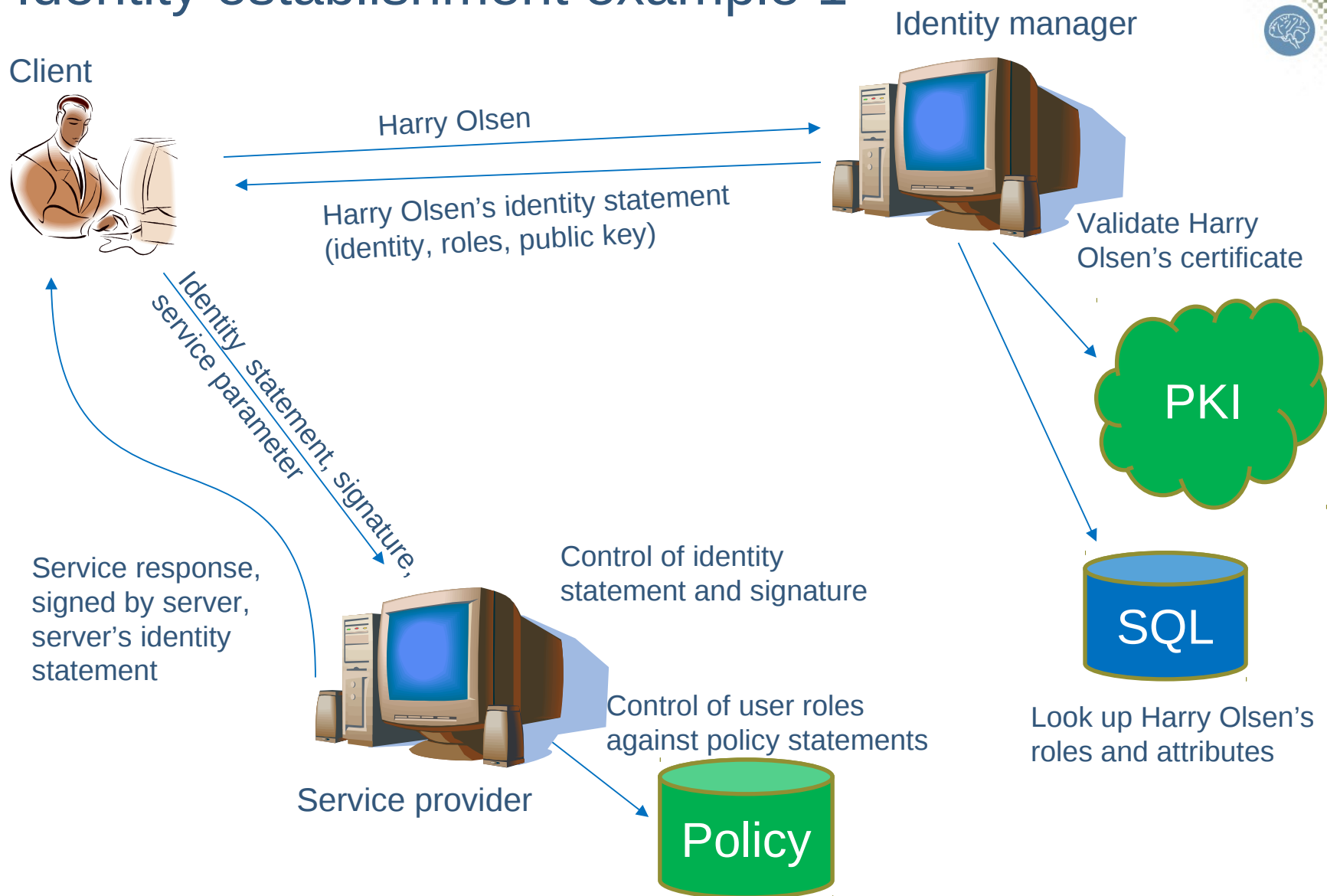
The Identity Statement

- Attested binding between properties and identifier
 - public key, attributes
- Signed by a trusted issuer
- Expires
- Both clients and services presents their identity statements in order to provide *mutual* authentication

Subject identifier
Subject public key/x509 cert
Subject attribute 1..n
Validity period (from-to)
Issuer identifier
Issuer public key/x509 cert
Issuer's signature



Identity establishment example 1





Trust assumptions

- The identity statement is issued (and signed) by the IdM
 - The service providers need trust in the IdM
 - that the identity statement are "correct"
- The service providers trust the authenticity of a client who demonstrates a private key (proof-of-possession)



Cross domain IdM principles:

- Inside each domain:
 - User key/certificate management
 - User roles/privileges management
- Between domains:
 - Trust in others' authentication process
 - Trust in integrity of user attributes
 - No management of foreign users
- Role based authorization process
 - since identity of guests are "unmanaged" in host domain

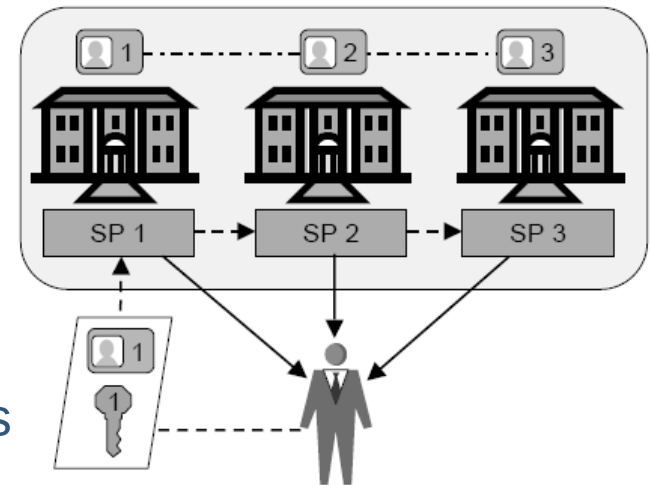
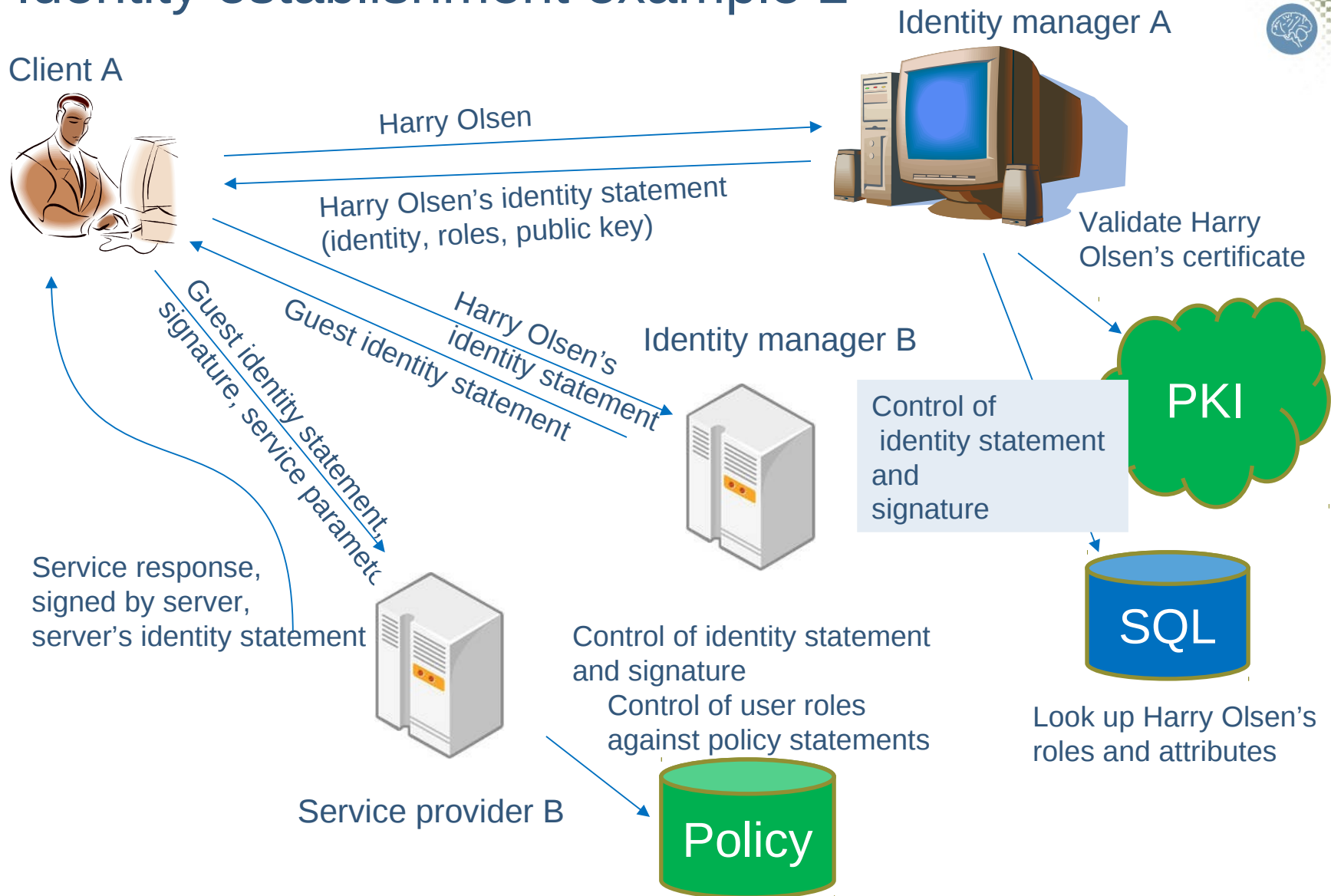


Figure 3: Federated user identity model.

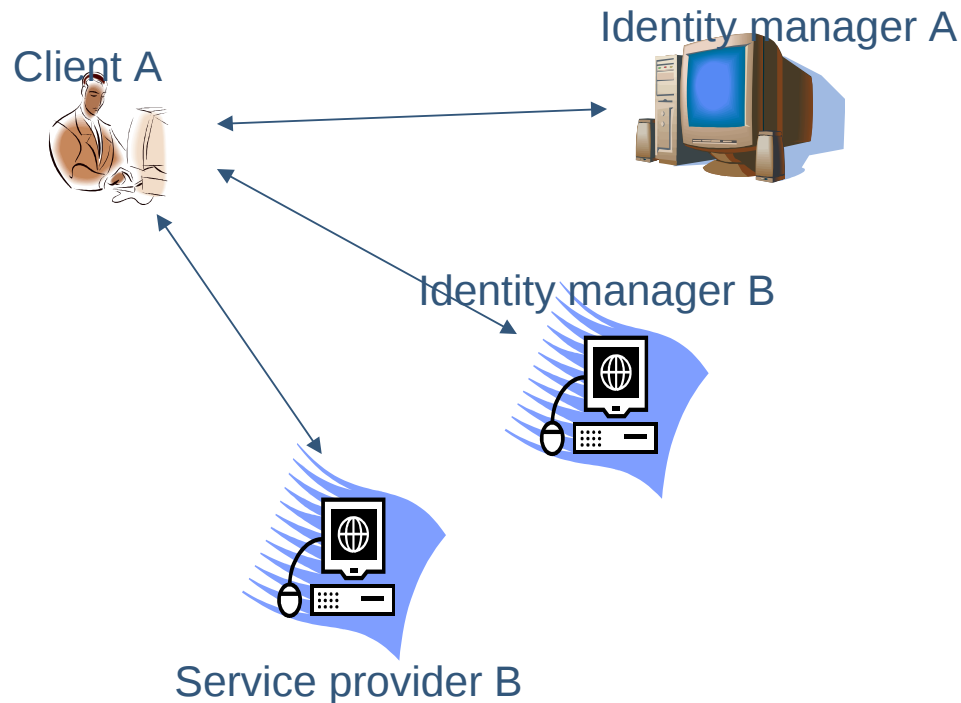


Identity establishment example 2



Trust relations

- IdM-B (Identity manager in domain B) trusts the authentication process of IdM-A.
 - it vouches for IdM-A by *re-signing* the identity statement
 - makes it into a domain-B security document





Caching is King

- The Identity Statement can be used for all servers in the domain for the duration of its lifetime
- Reduces traffic and connectivity requirements to the Identity Provider
 - Good for mobile systems



GISMO IdM's advantages

- Administrative and Authority Issues
 - autonomy of domains and COI
 - loose coupling between domains (certificate pair)
- Scalability issues
 - no CRL distribution
 - single domain user management
- Mobility / Tactical issues
 - occasional service invocations with IdM
 - client-A and server-B can connect independent on IdM reachability



Requirements for Cross Domain IdM:

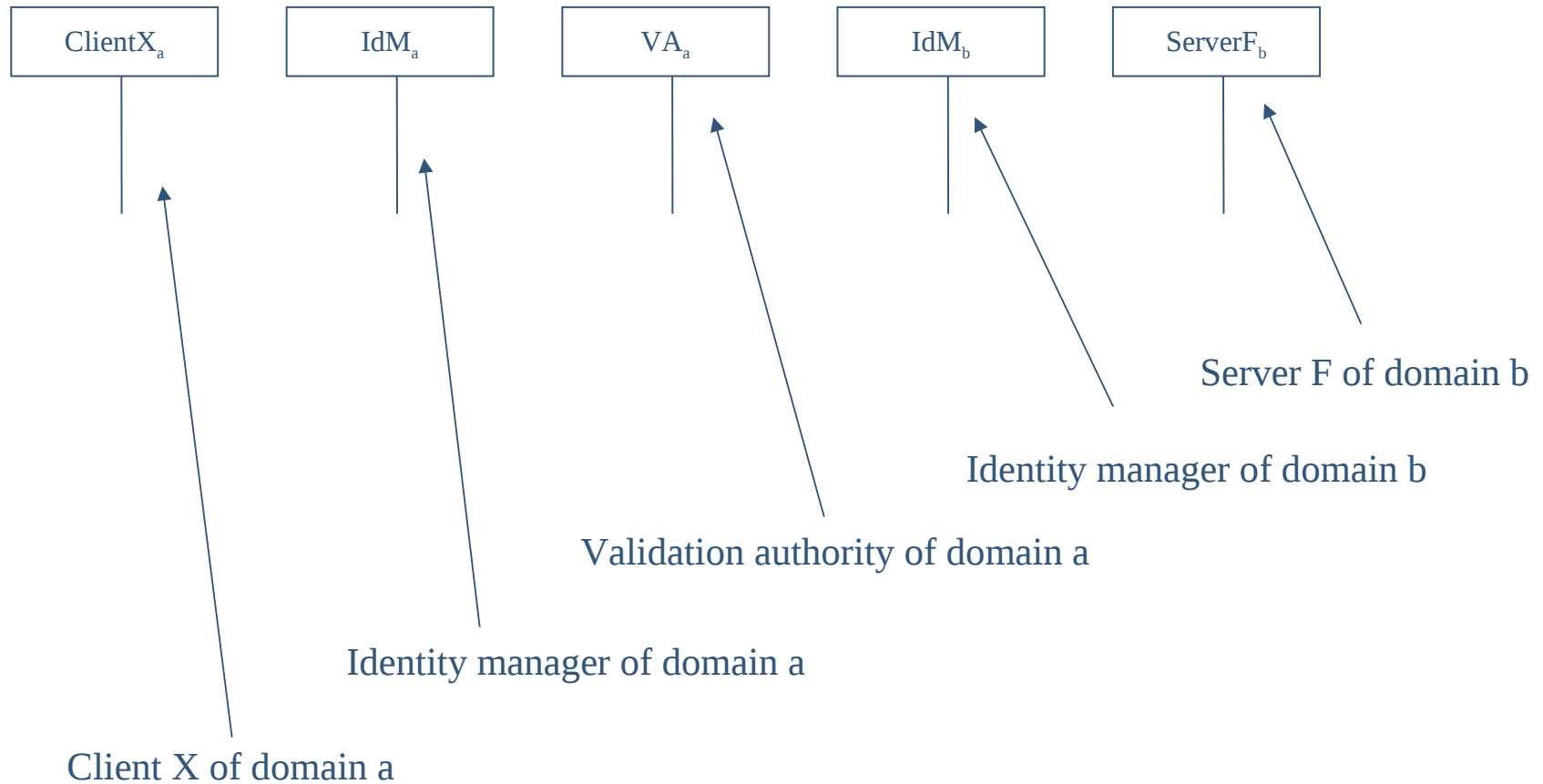
- Several security domains
- Low network traffic
- Few extra network invocations involved
- Uncomplicated and verifiable
- Suitable security level
- Application transparency



IS Issue and Authentication details

- In the following pages, a more detailed description of protocols and data structures is given.

The actors are:



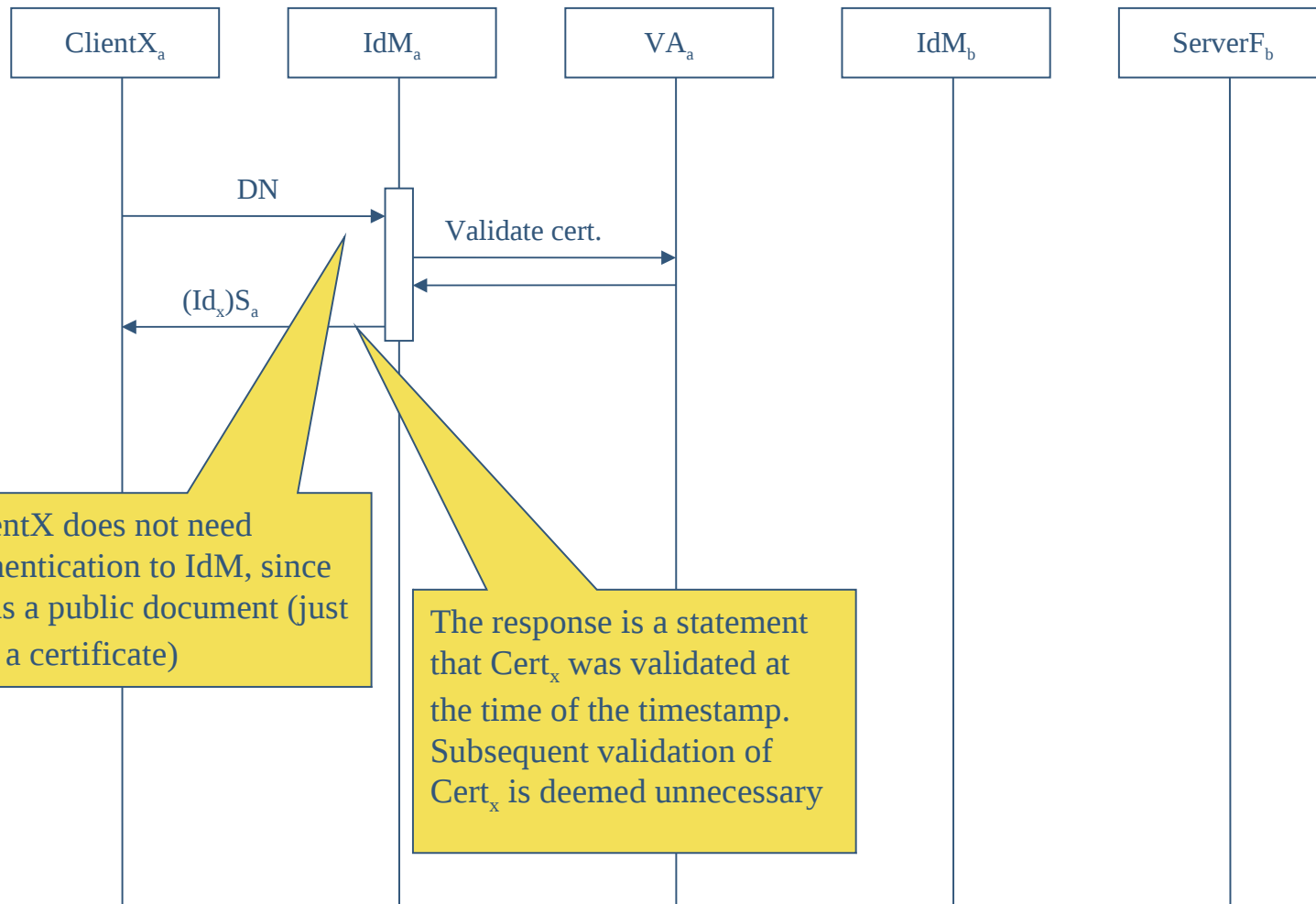


The data elements are:

- DN = X.500 Distinguished Name (CN=Anders Fongen, O=FFI..)
- Id = X509Cert + Timestamp + Attributes
- $(Id_x)S_a$ = Id of x signed by IdM_a
- $(Message)E_x$ = Message encrypted with public key of x
- $(Message)S_x$ = Message signed with private key of x
- Keep in mind:
 - Members of domain a (clients and servers) have trust in the signature of IdM_a
 - No revocation info is distributed, but handled inside the VA
 - The validity period of an Id is given by the timestamp

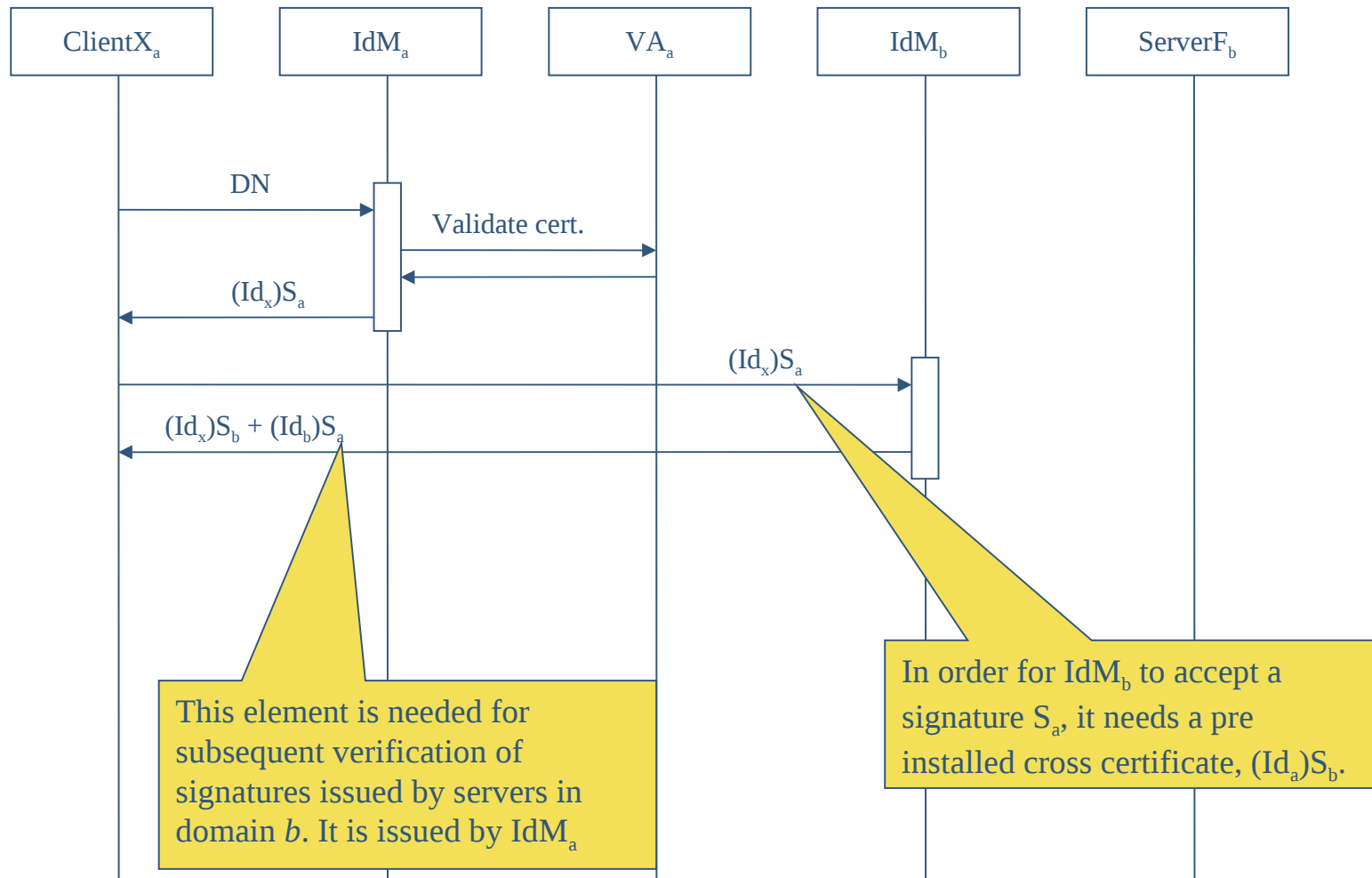


Step 1: Issue a signed Id (own domain)



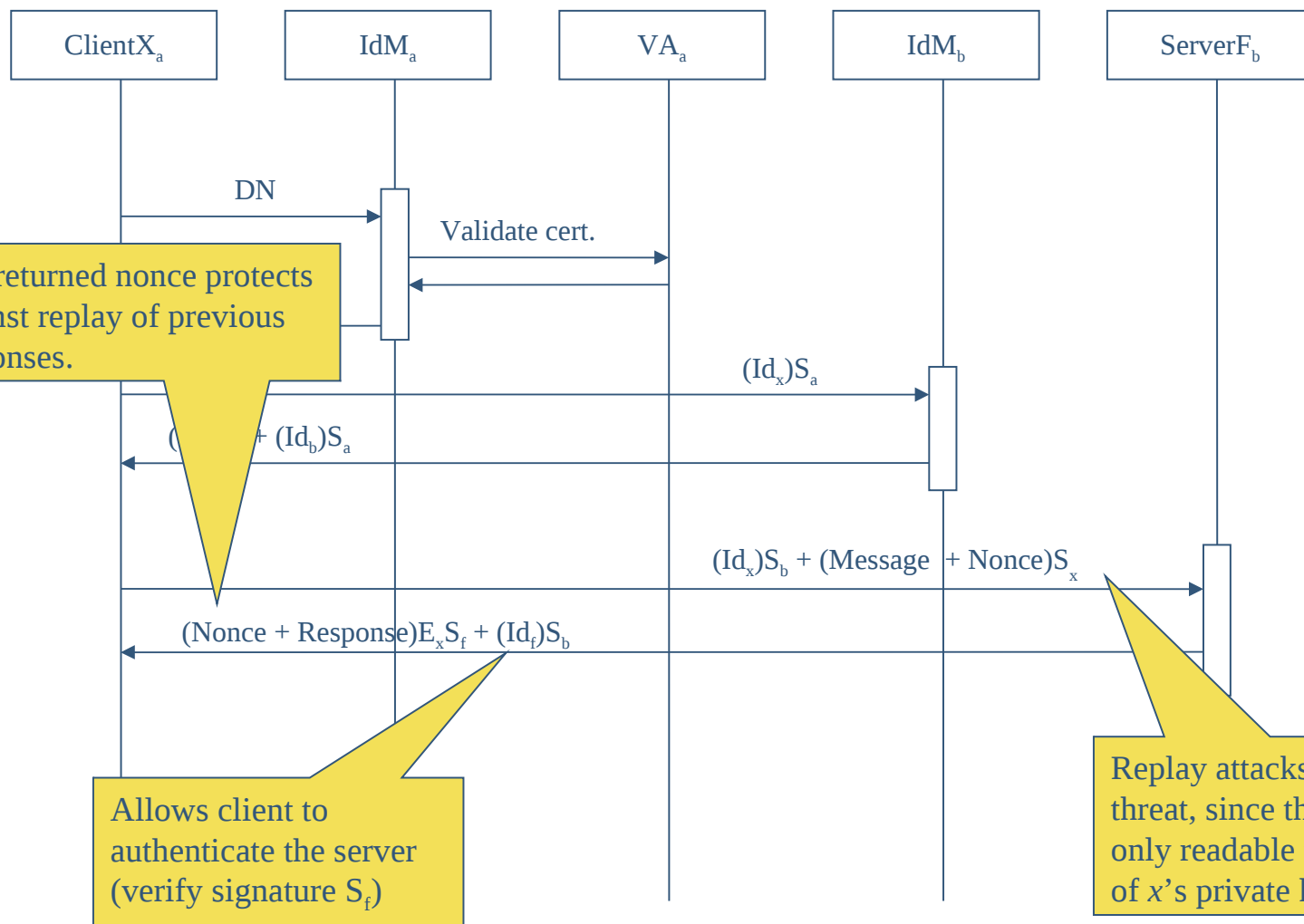


Step 2: Issue a guest Id (foreign domain)



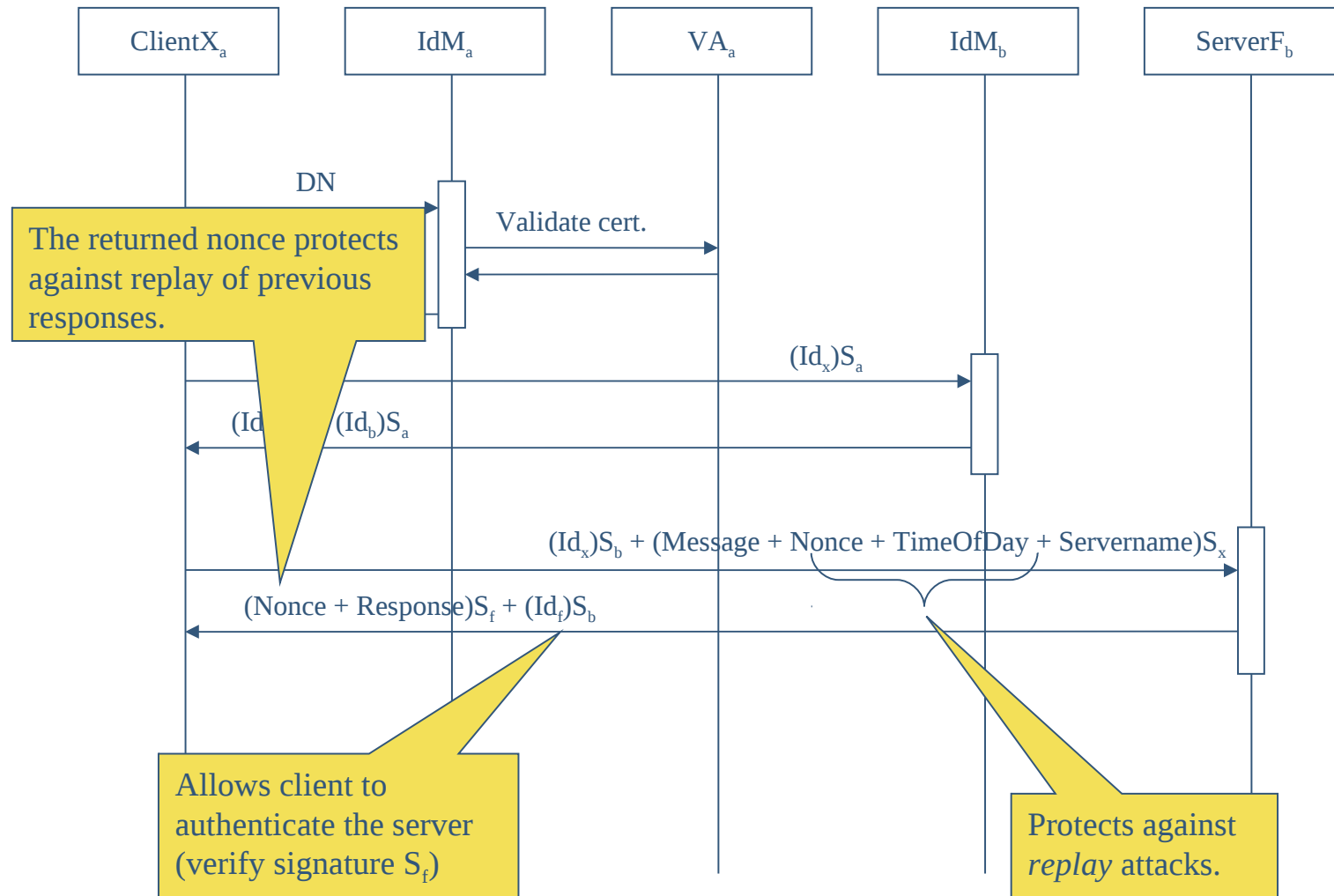


Step 3: Authenticated access to foreign server (stateless)

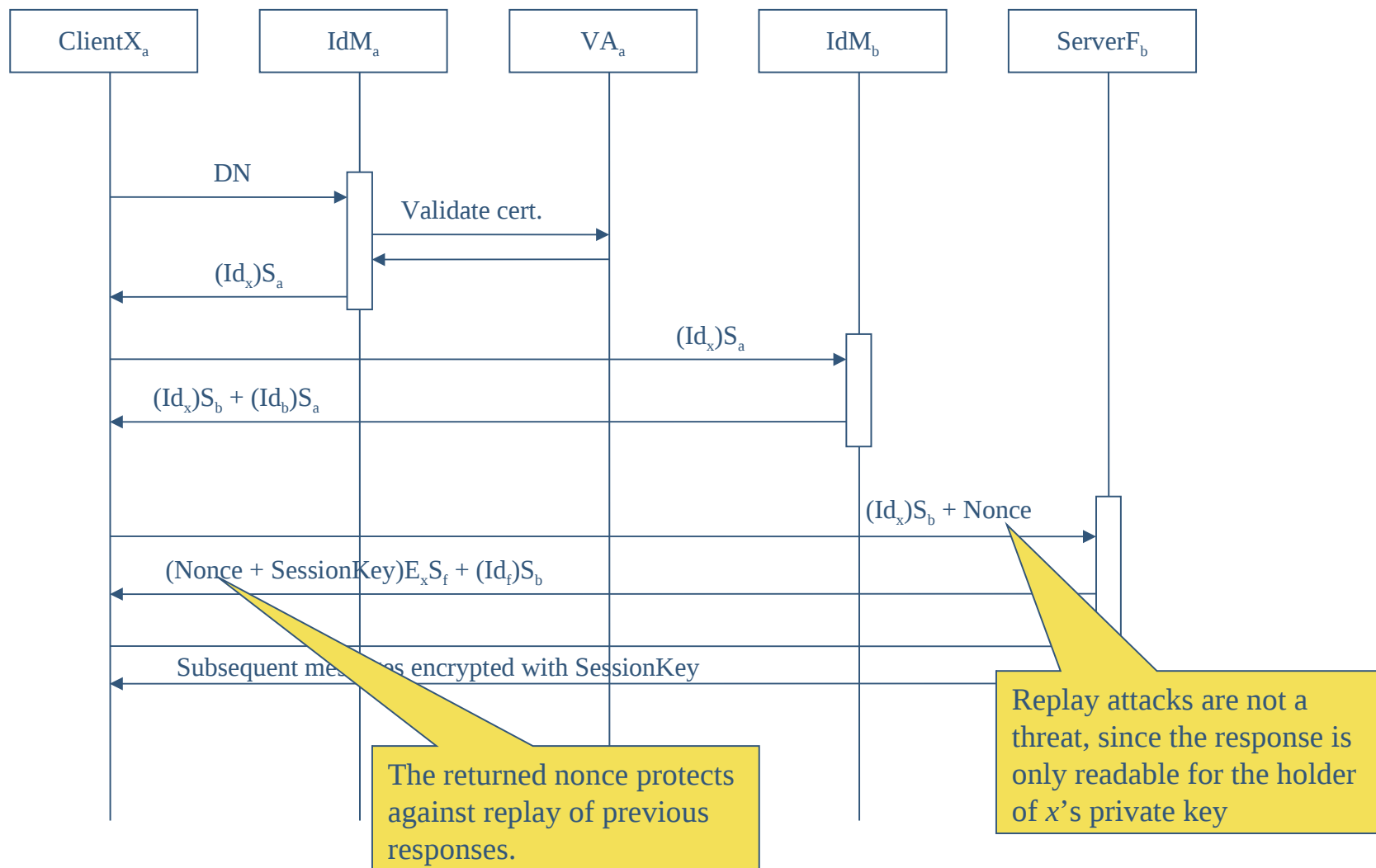




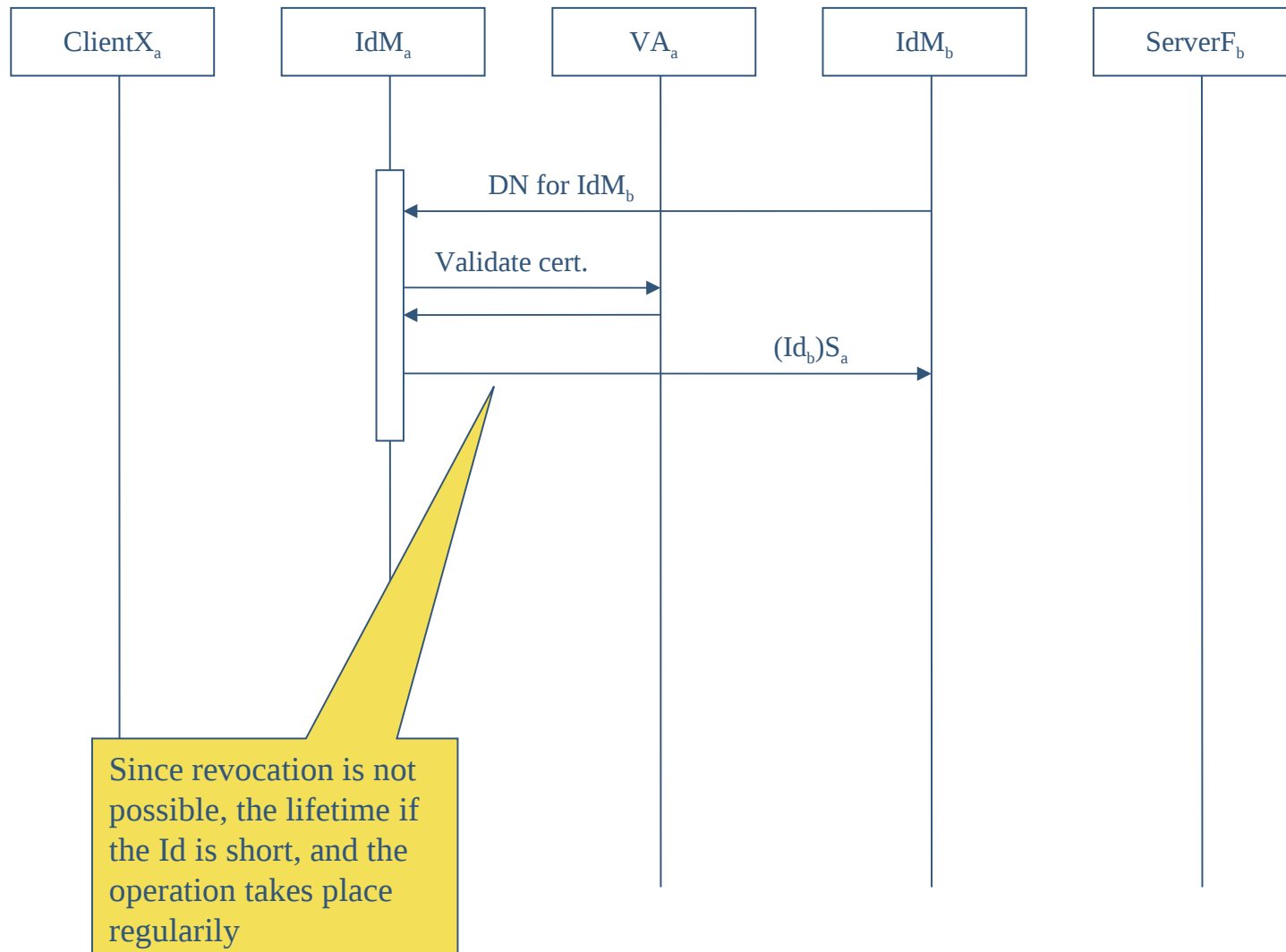
Step 4: Authenticated access to foreign server (stateful)



Step 5: Authenticated and encrypted access to foreign server (2-phase stateful)

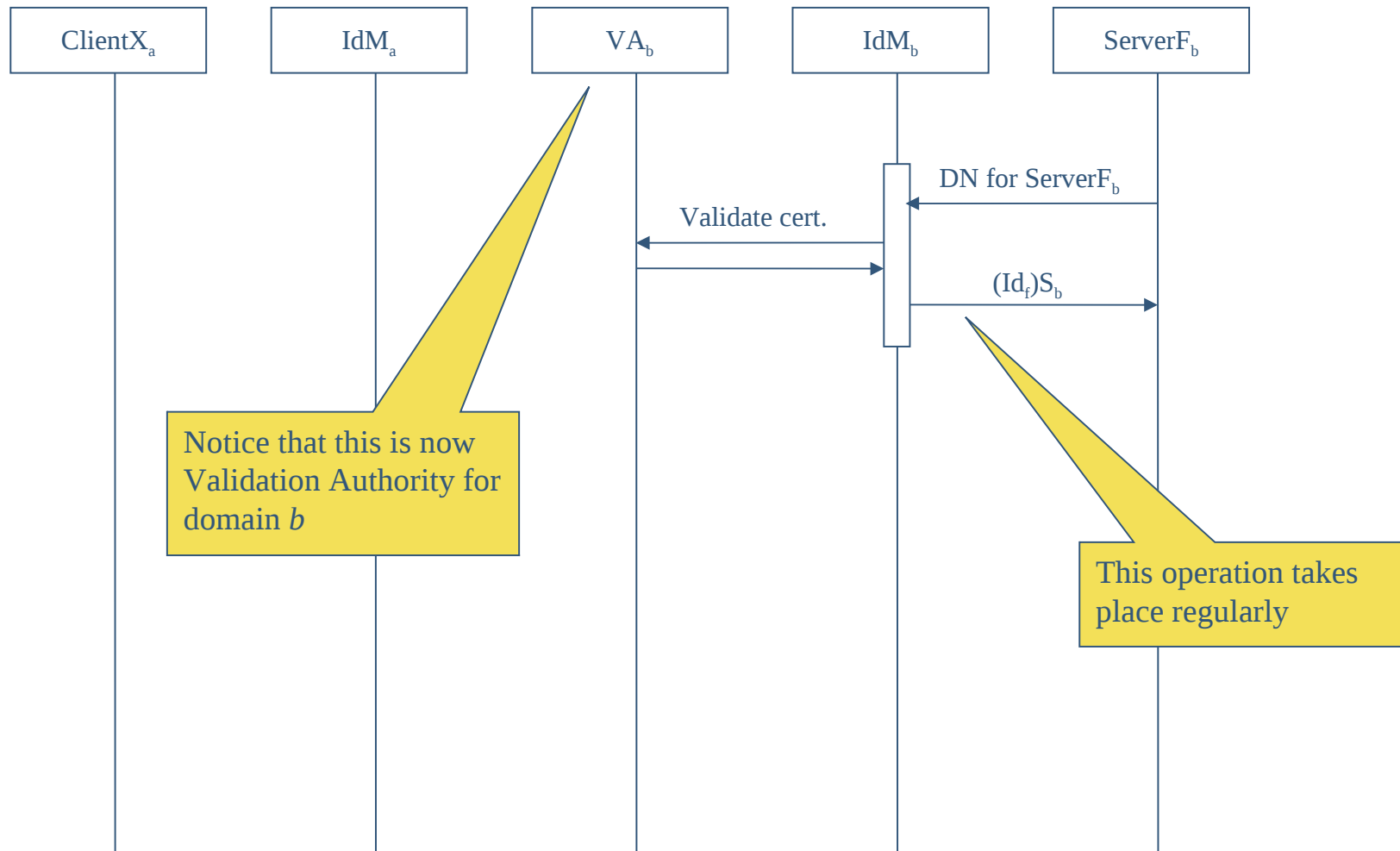


Side Step 1: Issue a cross domain IdM Id





Side Step 2: Issue a server Id





Scalability issues

- Administrative tasks relate to local users and resources only
 - scales with the size of own domain
 - **no CRLs distributed across domains!**
- Interdomain trust is expressed as cross certificates.
 - number of certificates are approx. domains^2
- Complexity of policies must be managed
 - roles harmonized across domain borders



GISMO IdM for mobile units

- Based on the existing GISMO IdM
- Extended to include support for Android units
 - With poor support for SOAP-derived standard like WSSec, SAML etc.
- Employs a dual stack presentation layer
 - SOAP and serialized Java objects (POJO)
- Interoperability issues are addressed
 - Proxy nodes breaks SOAP/POJO separation



The first solution was

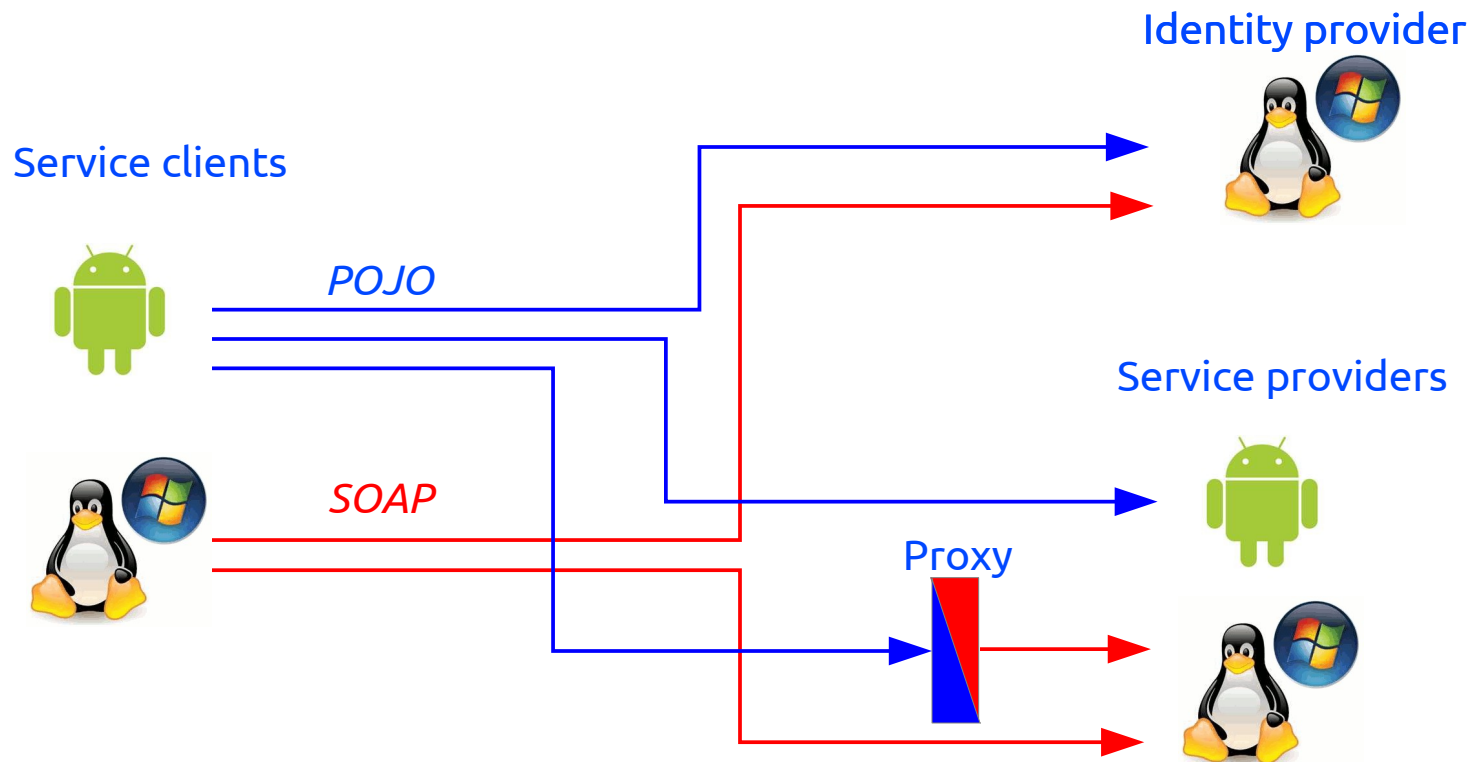
- Based on XML/WS standards
 - SAML, XACML, WS-Security
- Programmed in Java
 - Java-WS, Metro library
- Made as isolated/self-contained components
- *Not* portable to systems without extensive SOAP support.



GISMO IdM for Android

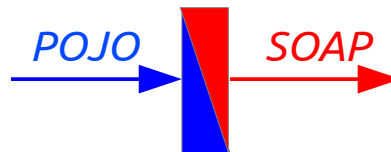
- Includes clients and service providers based on the Android OS
 - Android has poor SOAP support
 - No Sun XWSS library (WSsec and SAML)
 - No Web Services architecture
- Employs a dual presentation layer (protocol syntax)
 - SOAP/WSsec/SAML
 - Serialized Java objects (POJO)
- Offers a lightweight service container
 - Well suited for service provisioning on Android
 - Employs the *class property* of a Java object

Invocation diagram





Conversion proxy



Conforms to the service component definition

Parameter class selects the service

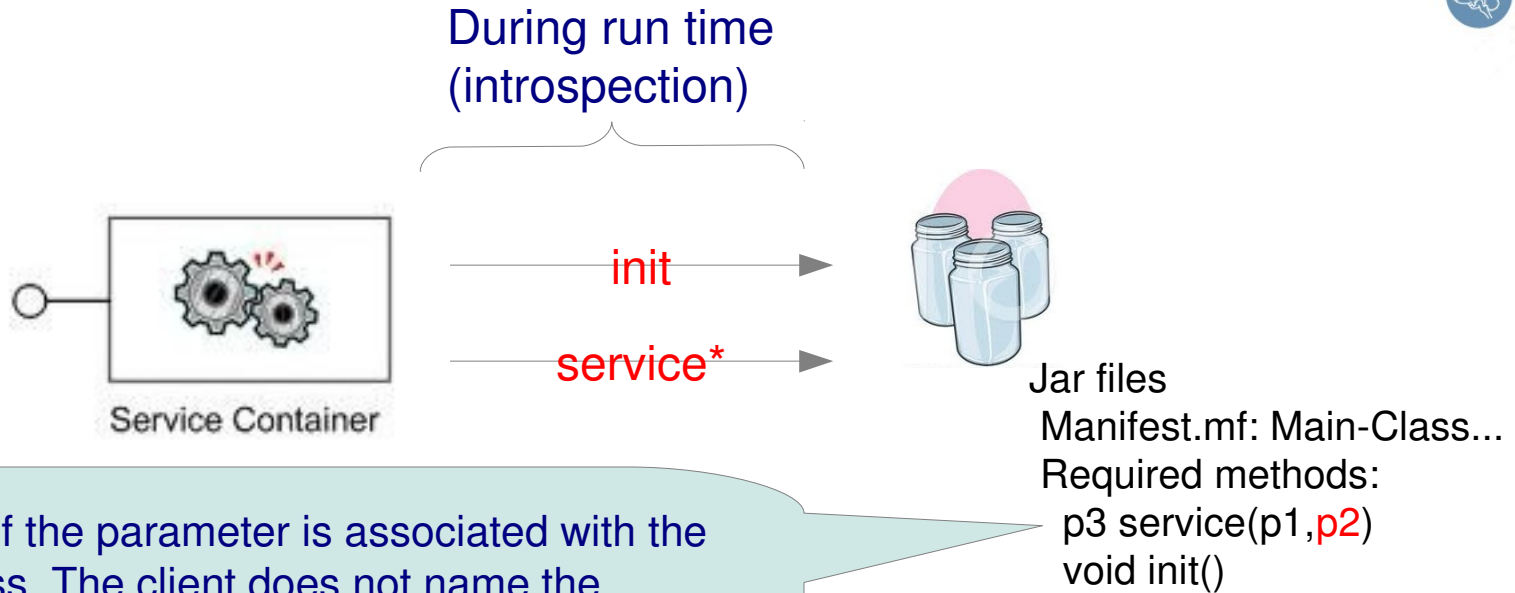
```
public class MainClass {  
    public Serializable service(WeatherRequest wr,  
                                Properties props) {  
  
        try {  
            Weather w = new Weather();  
            String result = w.getWeatherSoap()  
                .getWeather(wr.town);  
            return result;  
        } catch (Exception e) { return e; }  
    }  
}
```

Service class
generated from
WSDL-
compilation

NB: A conversion proxy
breaks the signature.



Service Components for POJO



CLIENT CODE:

```
try {
    ClientProxy client = new ClientProxy(prop);
    WeatherRequest p = new WeatherRequest(); // parameter
    p.town = "CHICAGO";
    Object s = client.invokeApplicationService(p);
    return s.toString();
} catch (Exception e) {...}
```

The light weight framework for service provisioning allows for Android units to offer services



Support for Access control

- Attested subject attributes can be used for
 - Roles in an RBAC context
 - Attributes in an ABAC context
 - Preferences for applications
- Maintenance of subject attributes are done in the IdM, not PKI
 - they reflect the policy of the COI, not PKI domain
- Cross-COI operation requires harmonization of attributes
 - names and values must mean the same everywhere
 - not necessary to exchange all attributes, only those significant for the cross-COI operation



Interoperability issues

- GISMO IdM only speaks «to itself»
- Turns into a *portability* issue
 - How can the protocols be implemented on other systems?
 - Present version written in Java
 - Ports well to Windows, Linux etc.
 - To the extent allowed by the libraries
 - Port to Python, .NET etc.
 - In «principle possible», but yet to be done
 - Relies on the availability of SOAP libraries (unknown)
- *Claim: POJO- and SOAP-based protocols are equally portable*

Experimental evaluation

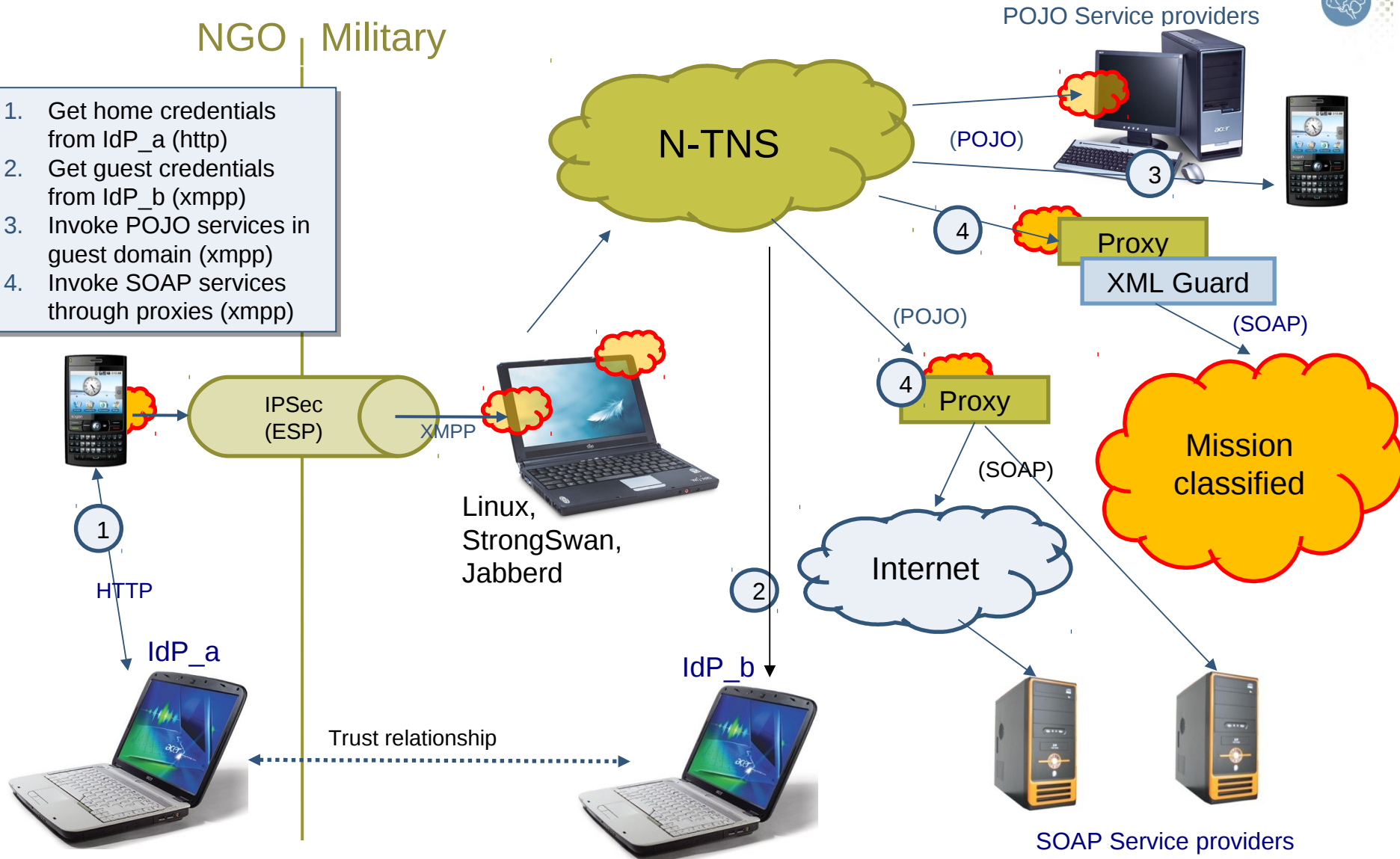
- Gismo IdM has been evaluated for correctness and feasibility
- A large field experiment will take place in June 2012
 - Opportunity to test stability, performance and scalability



Protected service invocations for Android

NGO | Military

1. Get home credentials from IdP_a (http)
2. Get guest credentials from IdP_b (xmpp)
3. Invoke POJO services in guest domain (xmpp)
4. Invoke SOAP services through proxies (xmpp)



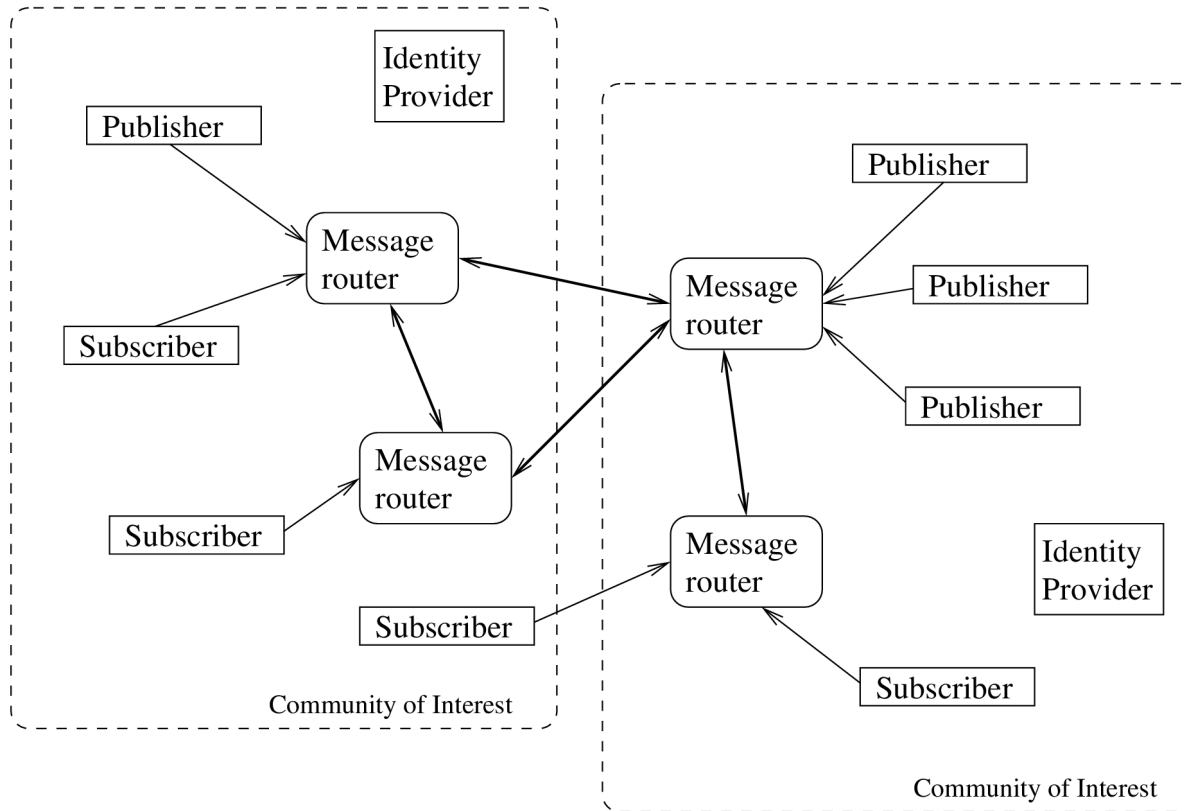
Principles of PubSub

- Messages are not individually addressed
- Routed from producer to receivers based on content
 - aka content routing
 - content metadata (topics)
- Publications are annotated with *topics*
- Receivers *subscribe* to topics
- Message transport is *asynchronous*
- Scales better
 - employs multicast topologies
- Smaller resource consumption
 - Messages are queued outside the hosts

PubSub security problems

- Is the published information authentic and unmodified?
- Who is the publisher, is it authorized?
- Will the published information only reach authorized subscribers?
- Who operates the message routers? Are they operating correctly?

Extending the PubSub network with Identity Management



Principles for PubSub security

- The confidentiality is the concern of the *Publisher*
 - should be able to express *subscriber requirements*
- The integrity is the concern of the *Subscriber*
 - should be able to express *publisher requirements*
- The authenticity of identities is the responsibility of the *Message Routers*
 - through identification of clients and peer message routers
- The correctness of keys and credentials is the responsibility of the *Authority*
 - establishes identities, issues and revokes keys and credentials

How to express a subscription?

- One or more topics
 - must match (hierarchically) the topics of the publication
- A publication requirement
 - a boolean expression evaluated with respect to the publisher's identity attributes
 - must evaluate to *true*

Subscription

topics:

/news/security/incidents

/alarms/ids/building-203

publisher requirements:

(nationality="uk") or

(sec_clearance="NATO secret")



IdM and Integrity Control

- Trust in security operations relies on service integrity
 - unadulterated software and hardware
- My means of hardware, a sterile configuration can be *sealed*
 - checked at bootstrap or anytime

$$hmac = f(K, h(mem), challenge)$$

- attested by an IdP and validated by anyone
- Can relax other security requirements

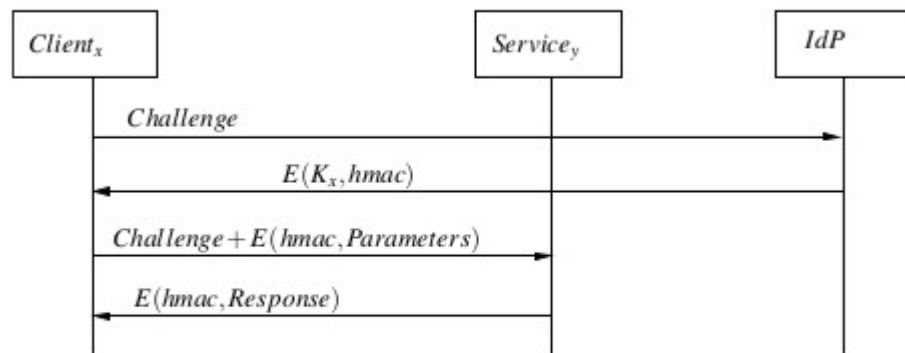


Figure 1. The protocol for protection of genuineness, based on symmetric crypto keys and the HMAC function given in Equation 1.



Publication from these efforts

- *“Identity Management Without Revocation”*, Securware 2010, Venice, Italy, July 2010
- *“Architecture Patterns for a Ubiquitous Identity Management System”*, under review for ICONS 2011, St. Maartens, January 2011.
- *“Identity Management for Android”*, Securware 2011
- *“Federated Identity Management in a Tactical Multi-Domain Network”*, Int. Journal on Advances in Systems and Measurement, 2011 vol 4 nr 3&4
- *“Identity Management and Integrity Protection in Publish-Subscribe systems”*, under review for MILCOM 2012
- *“Identity Management and Integrity Protection in the Internet of Things”*, under review for EST 2012
- *“Protected and Controlled Communication Between Military and Civilian Networks”*, under review for MCC 2012



Summary and Conclusions

- An IdM should support authentication as well as access control
- An IdM should not revoke “live” identity statements
- An IdM MUST support mutual authentication
- An IdM should provide loose coupling between domains and maintain autonomy
- Android has poor support for SOAP, SAML and WSSec.
- A separate presentation layer can be defended
 - Serialized Java objects have been chosen
 - Flexible, efficient and intuitive API
- Invocation “across” presentation layers (POJO-SOAP) is provided by proxy agents
- Future research includes an international field experiment